

Securing Resource Constrained Processors with Name Confusion

Mohamed Tarek Ibn Ziad, Miguel A. Arroyo, Evgeny Manzhosov, Vasileios P. Kemerlis, and Simha Sethumadhavan



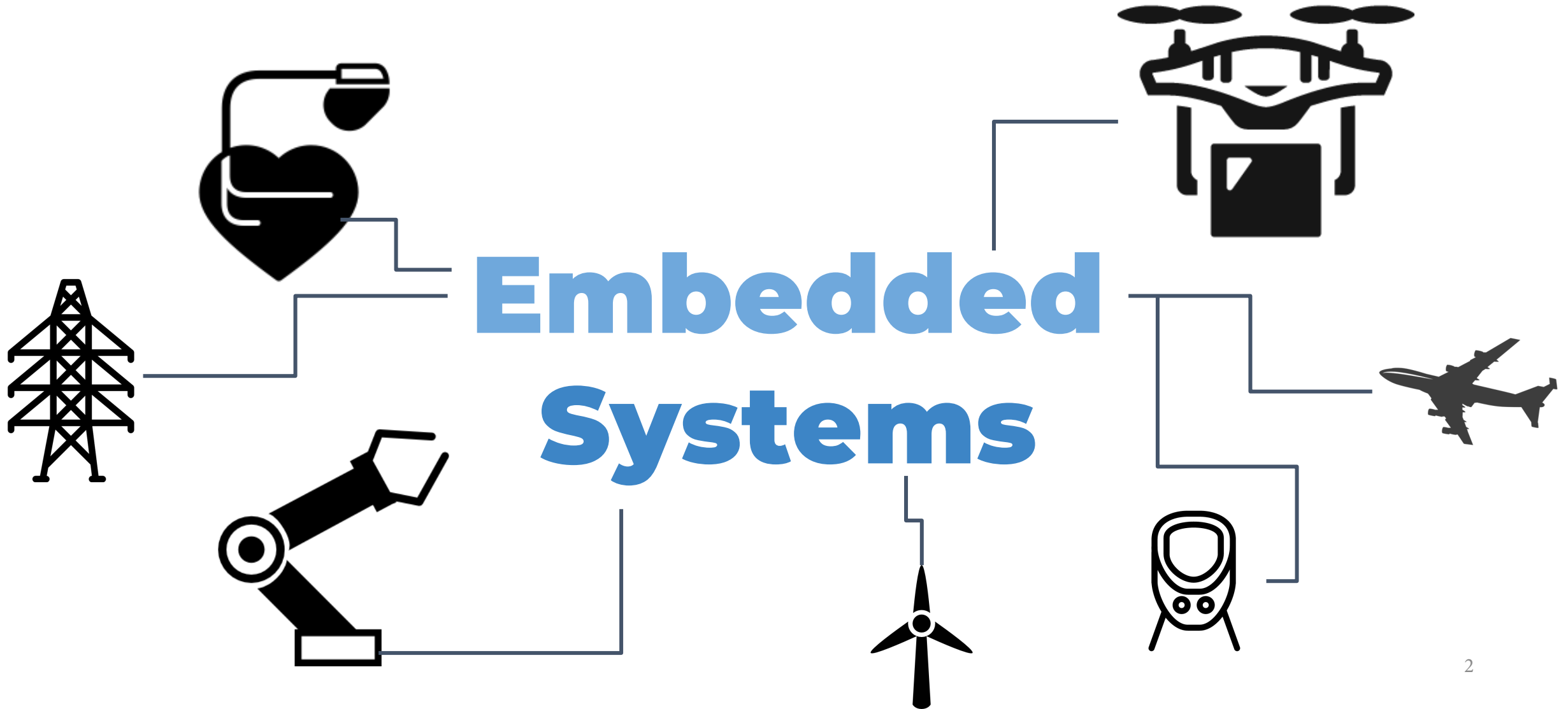
COMPUTER SCIENCE

Columbia University
Brown University
09/21/2021



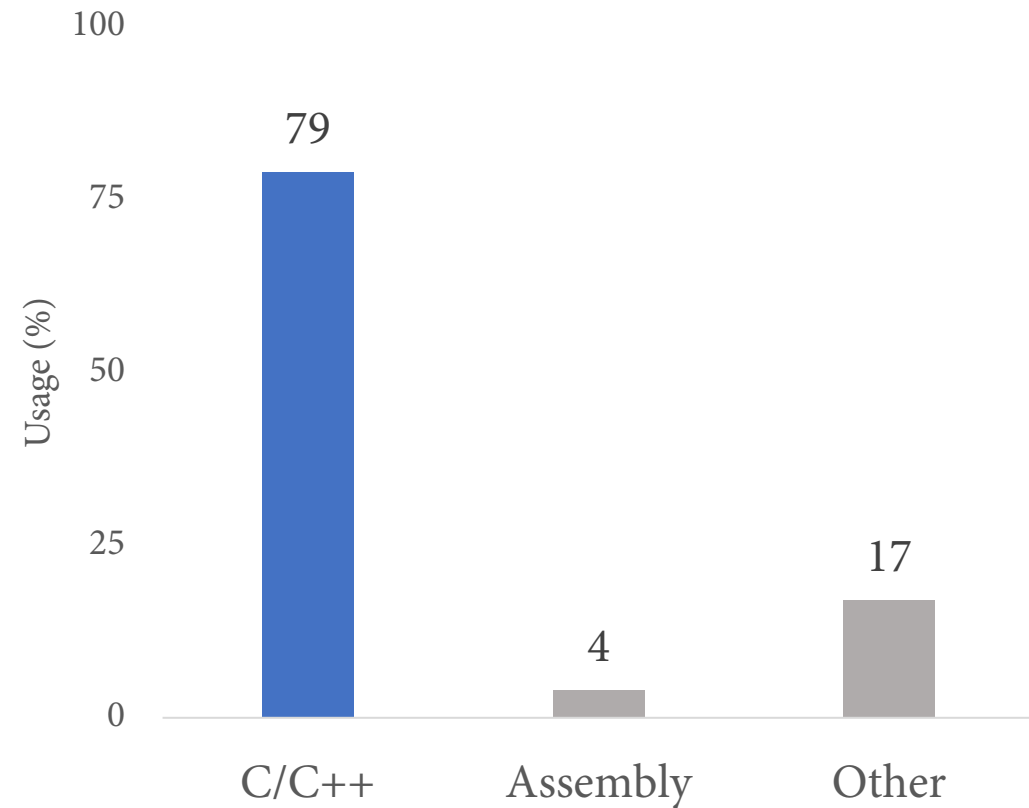
BROWN

Embedded systems are everywhere!

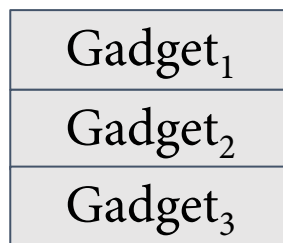


Why focus on software threats?

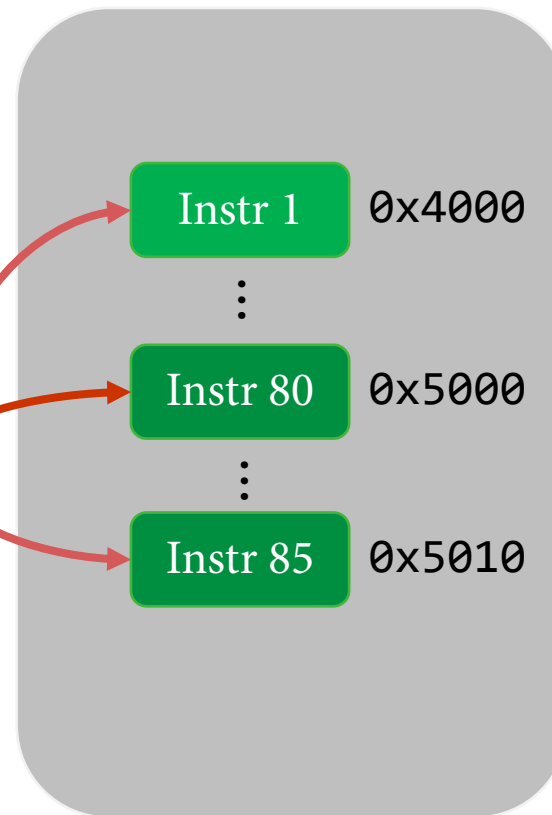
CPSs are predominantly written in memory unsafe languages.



Code Reuse Attacks



Stack

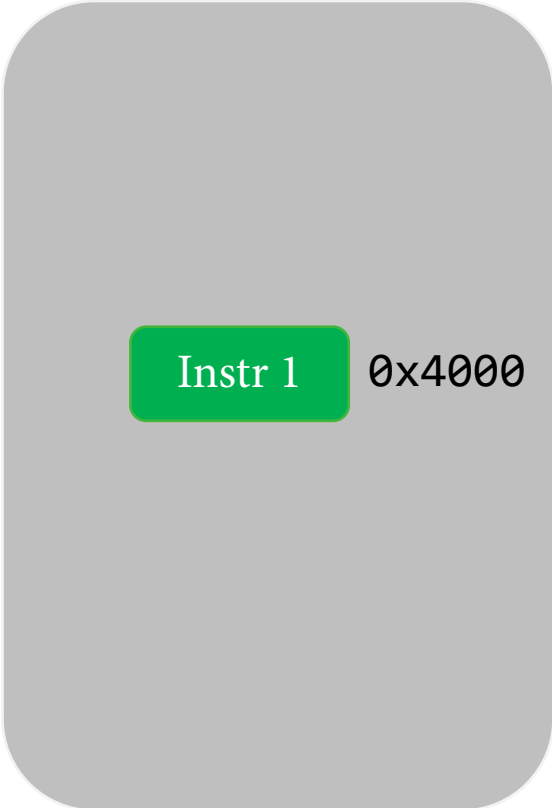


Virtual Address Space



Traditional Architecture

An instruction has a single name (or address).

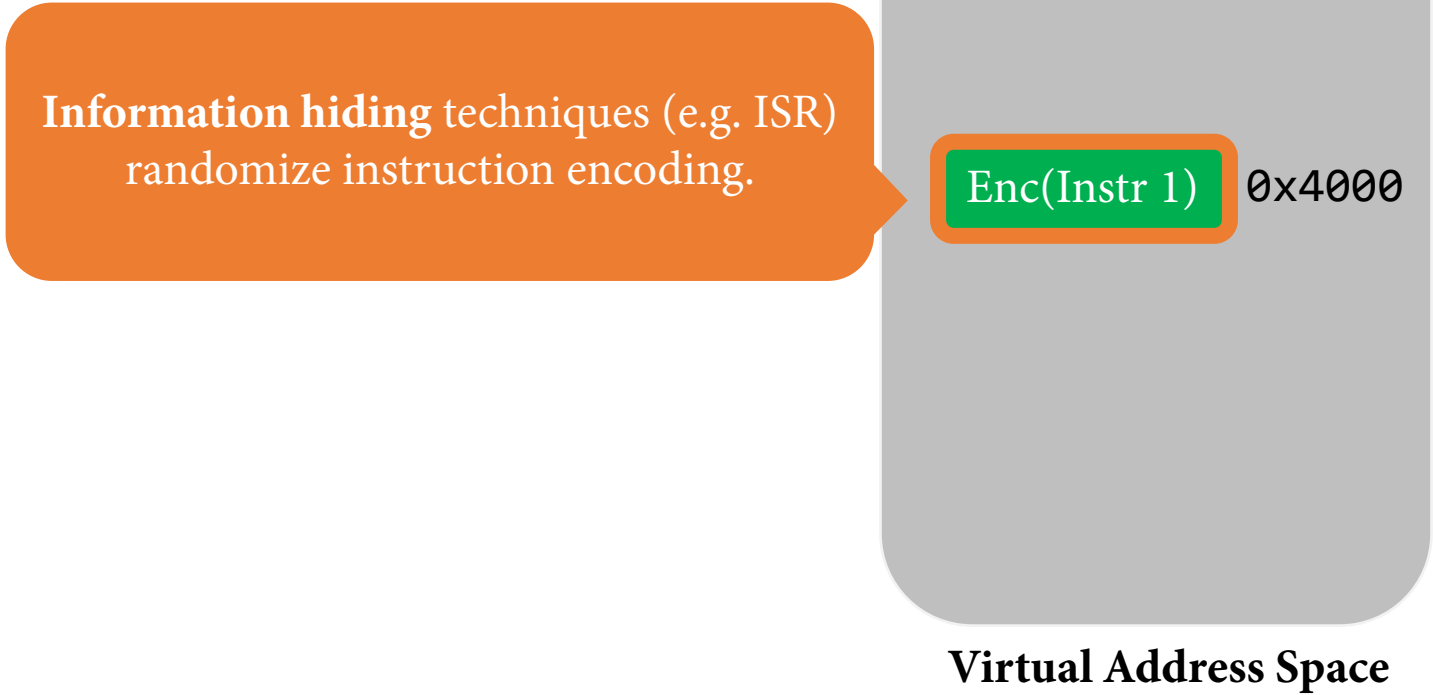


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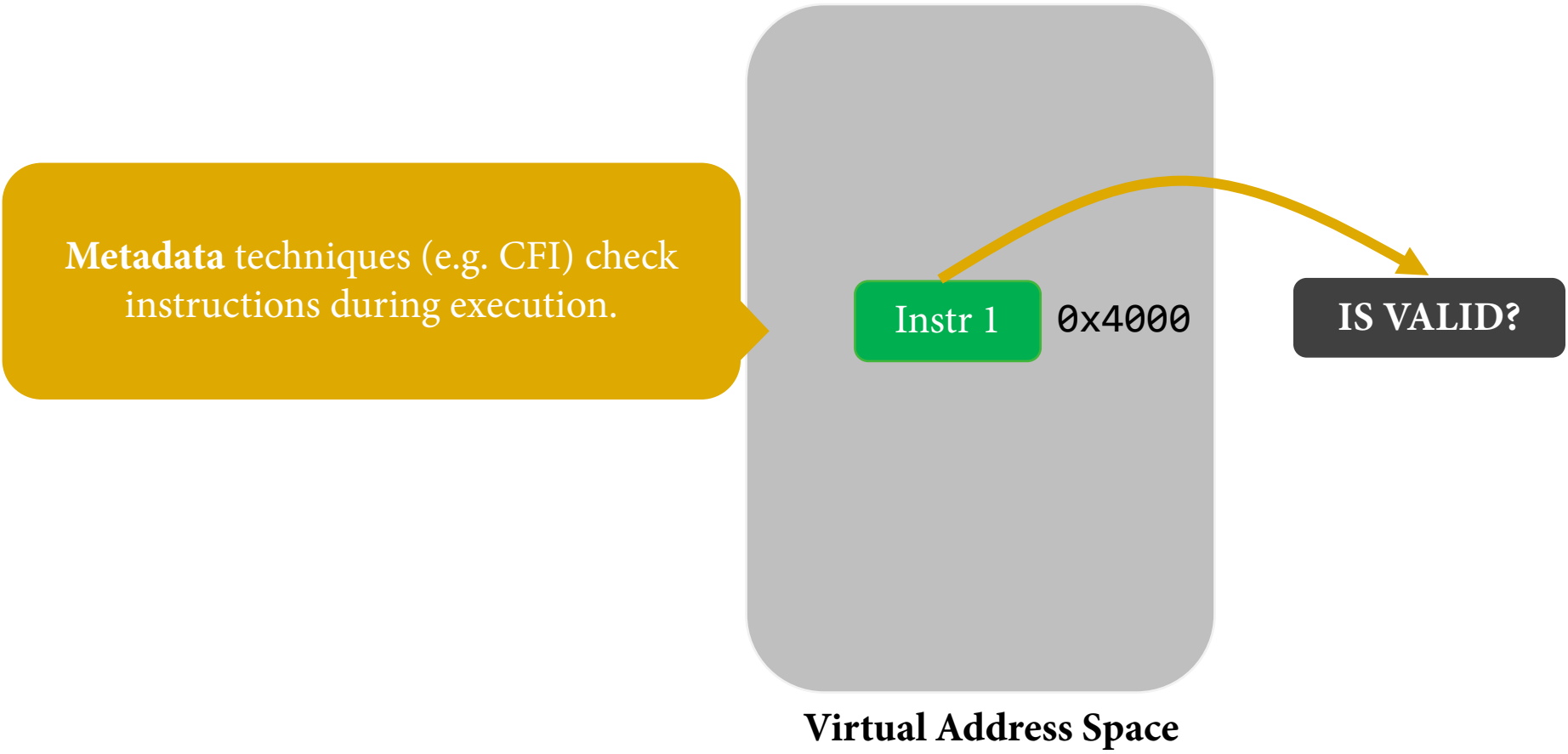
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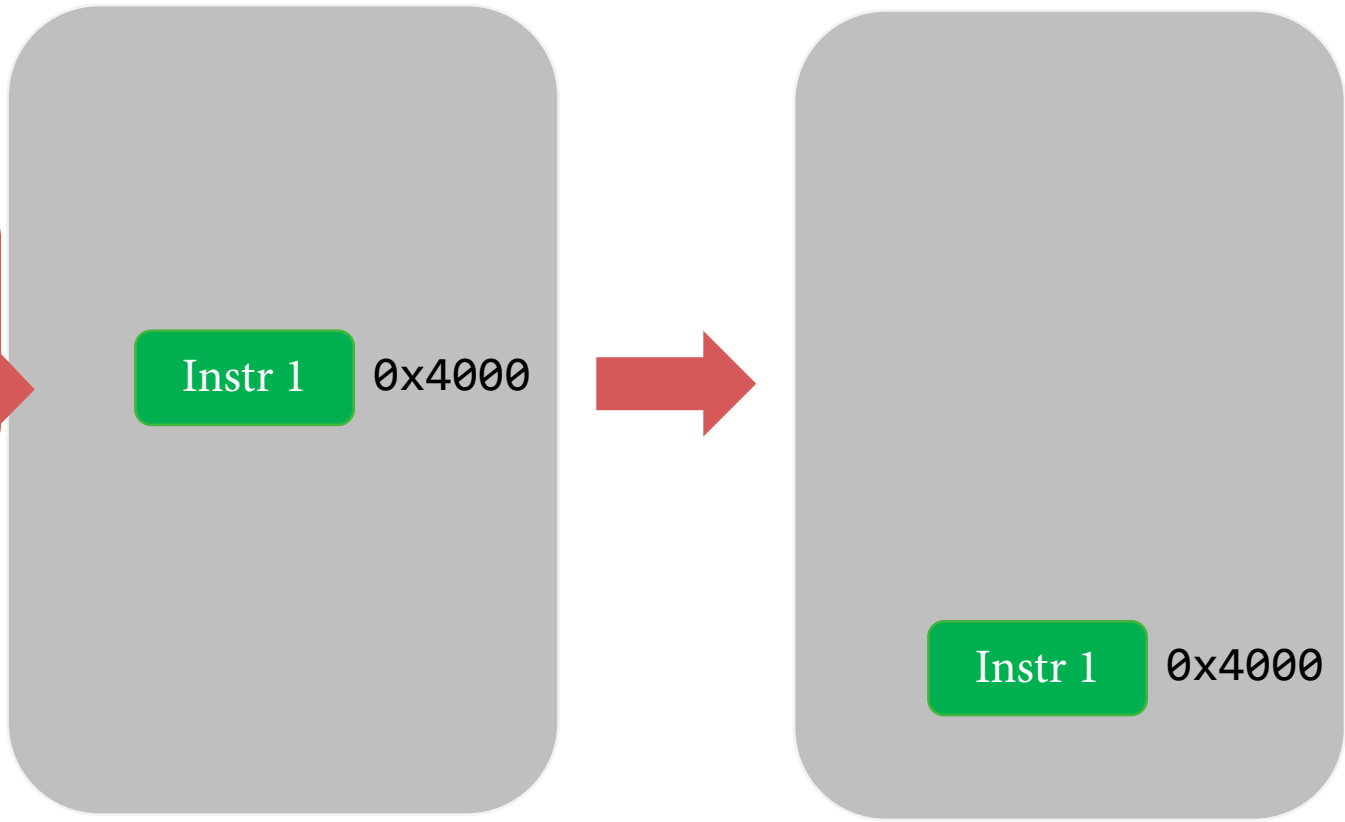




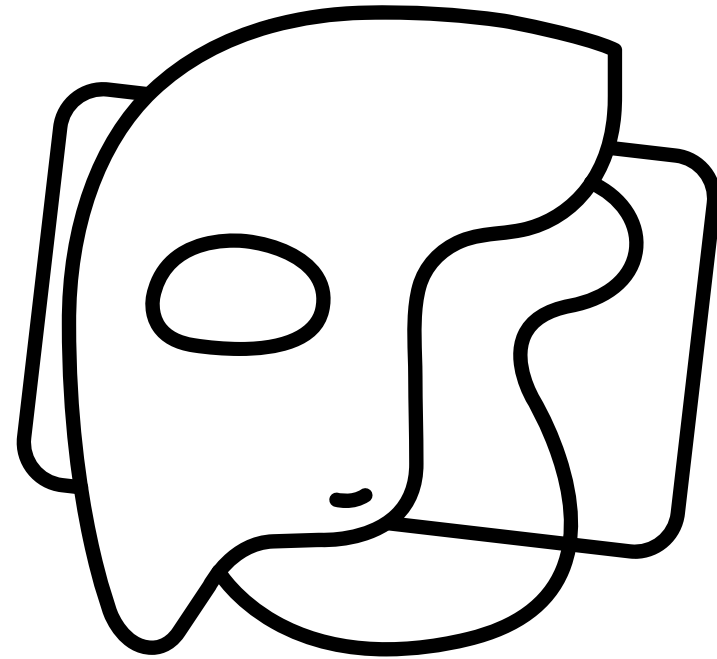
Traditional Architecture

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Moving target techniques change instruction names over time.



Virtual Address Space

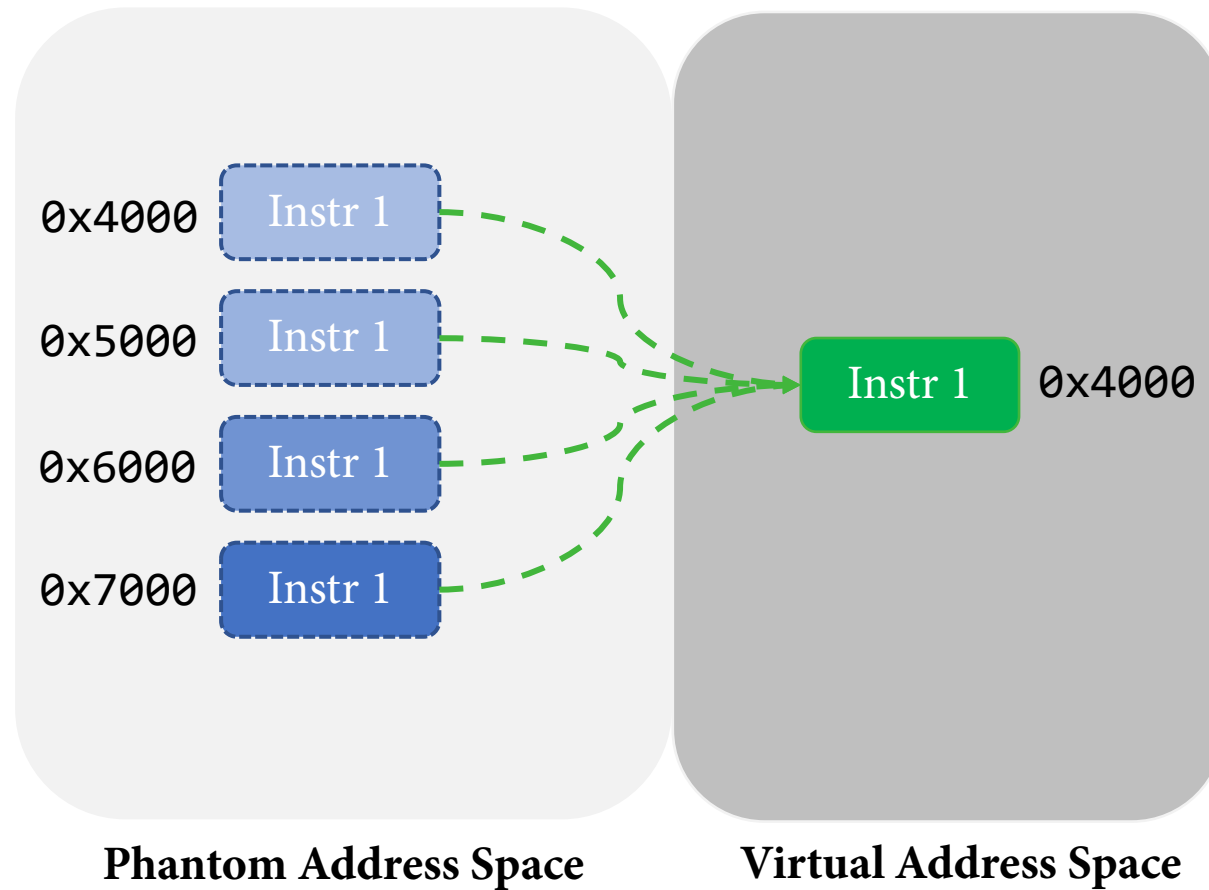


Phantom Name System

A Name Confusion Architecture

Name Confusion Architecture

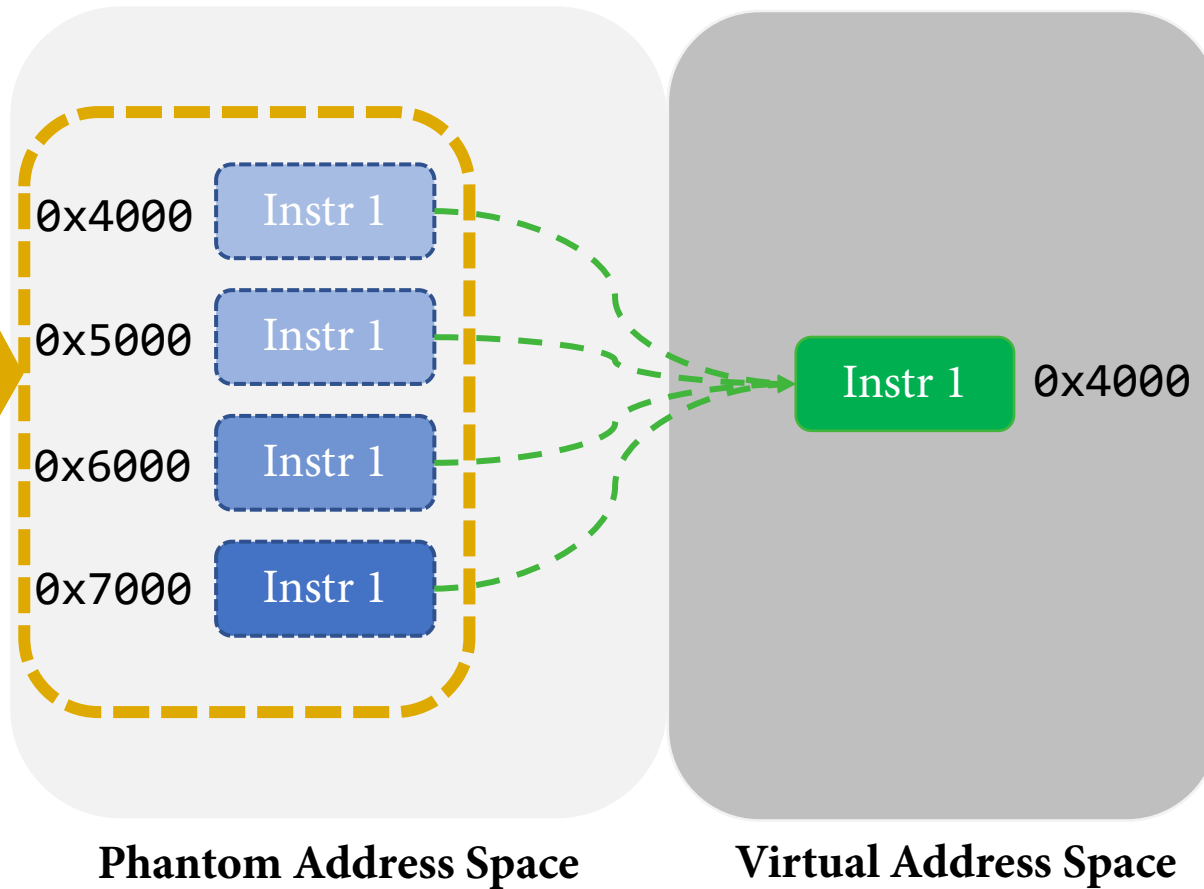
Multiple phantom addresses alias to an instruction.



Name Confusion Architecture

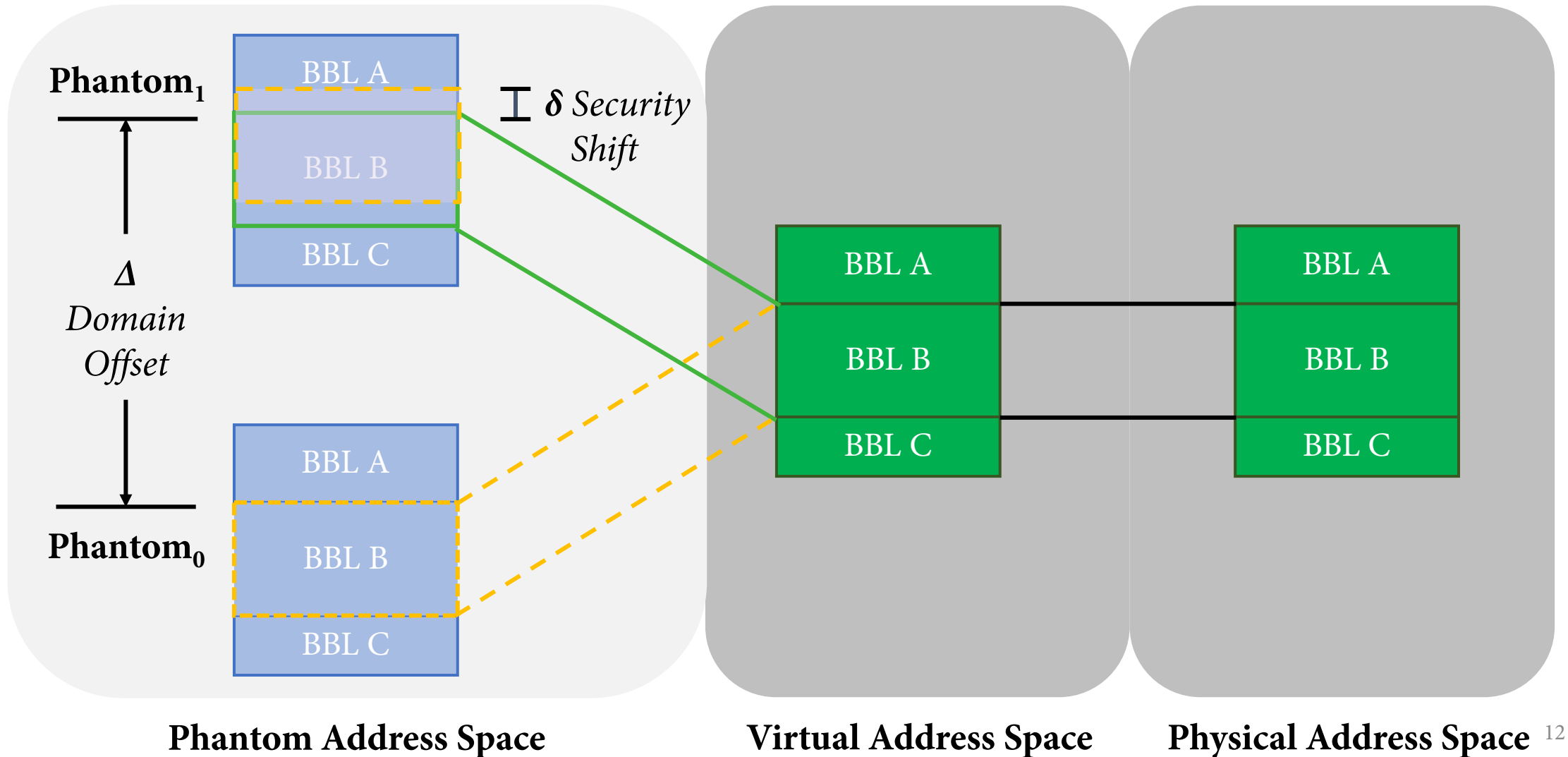
Multiple phantom addresses alias to an instruction.

An attacker must **guess** which will be executed!



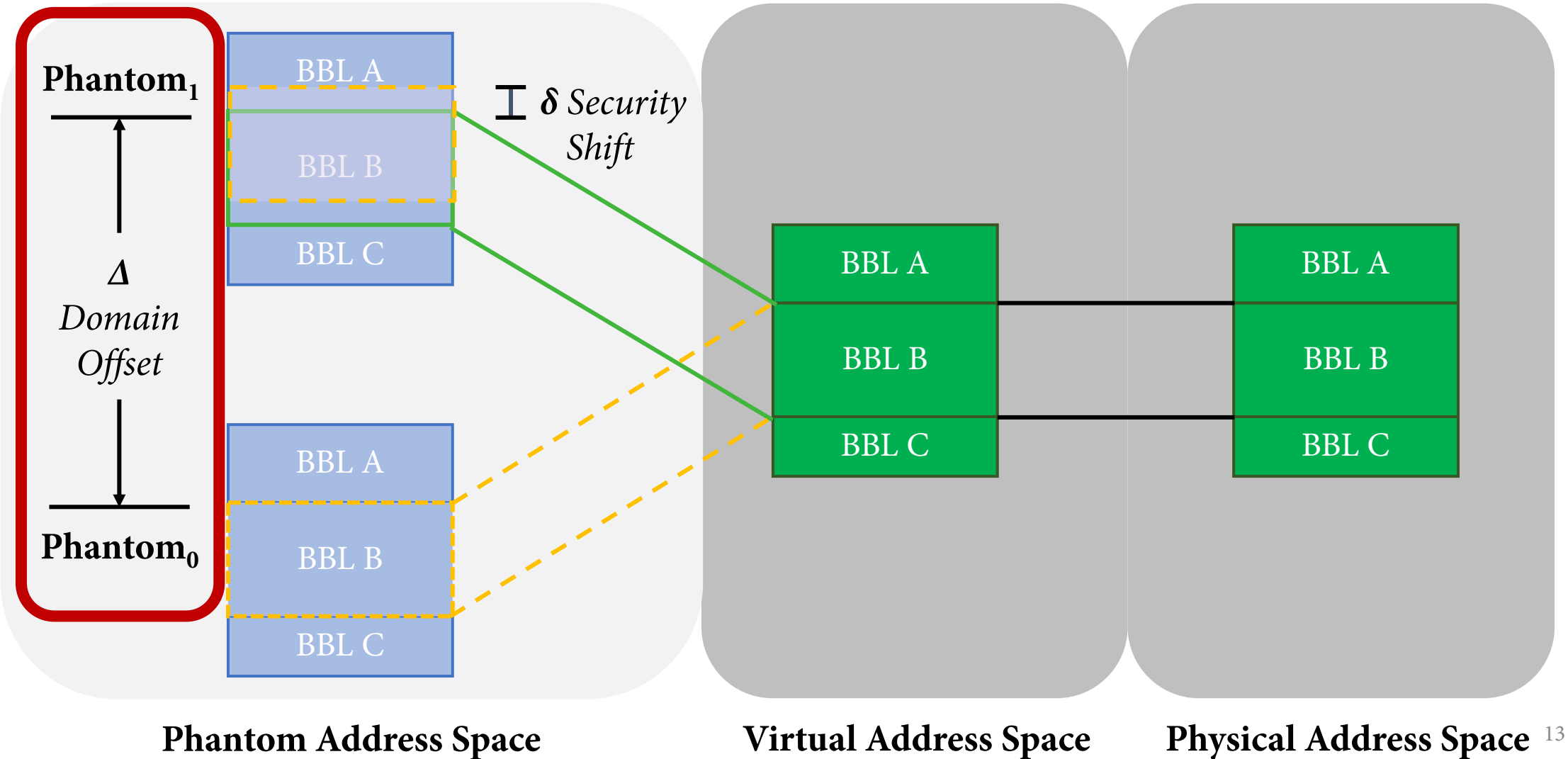
How are phantoms constructed?

Phantoms are logically displaced relative to the original program.



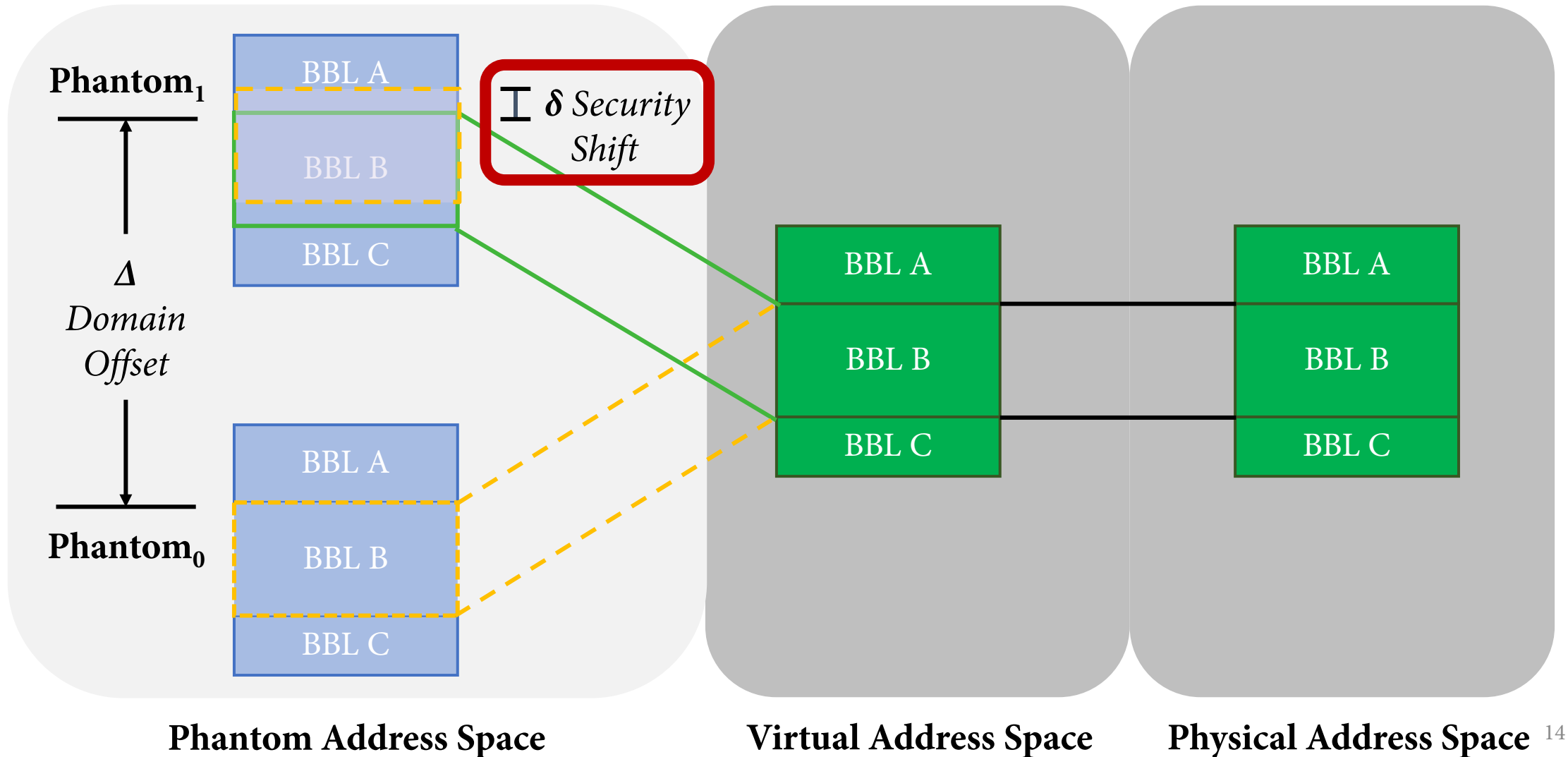
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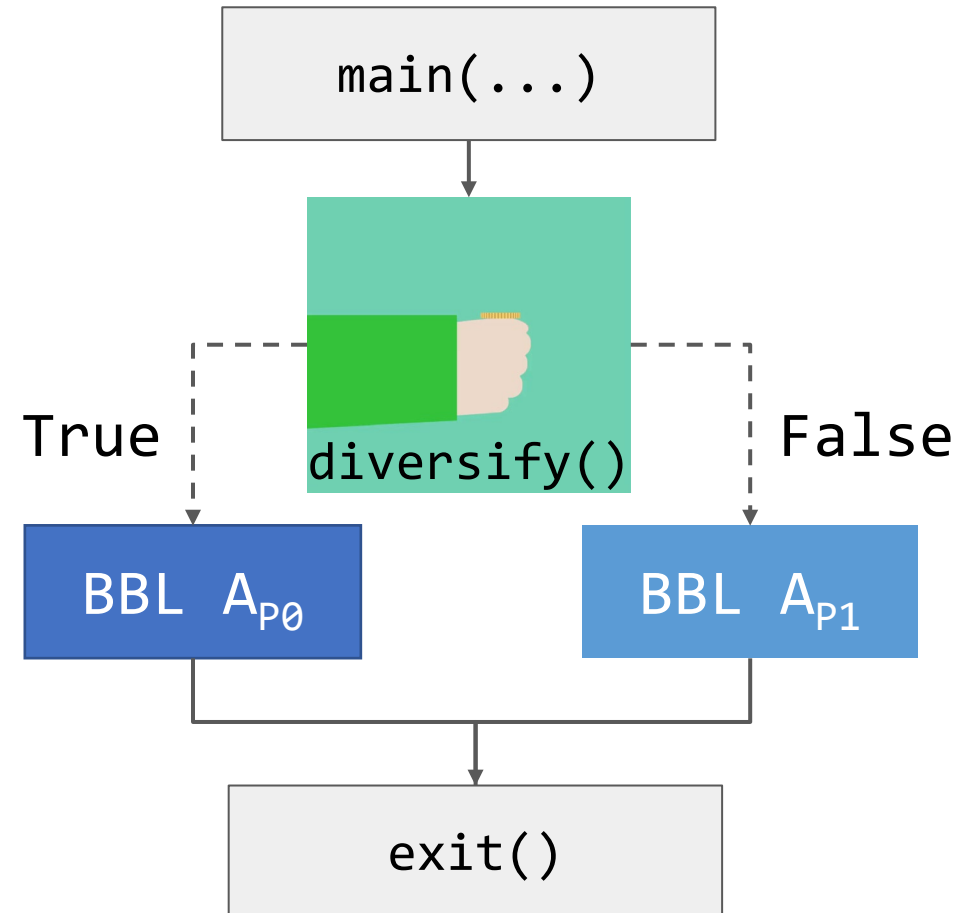
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How does PNS diversify execution?

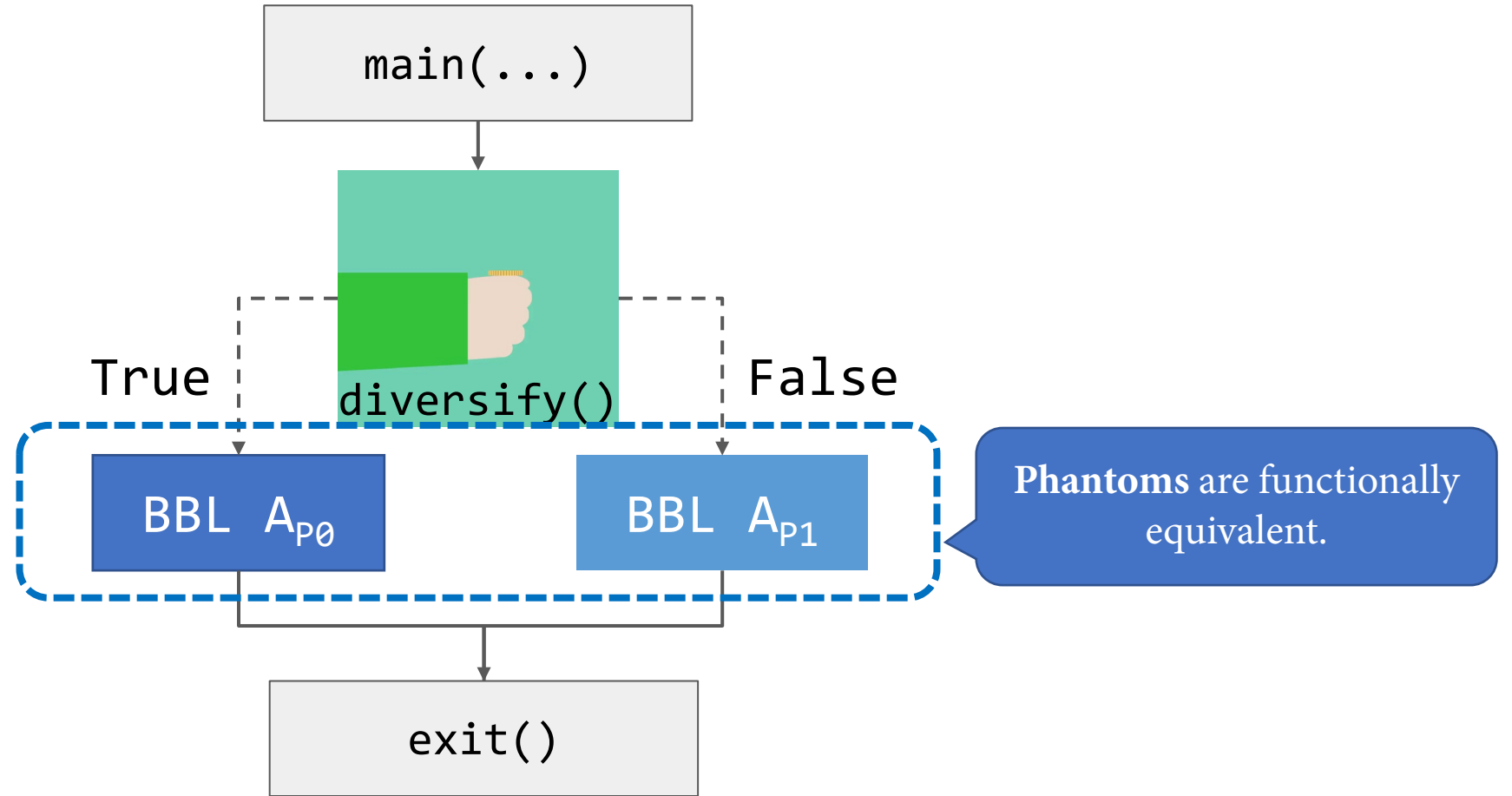
It diversifies the *path* of execution at every basic block.



Program Control Flow Graph

How does PNS diversify execution?

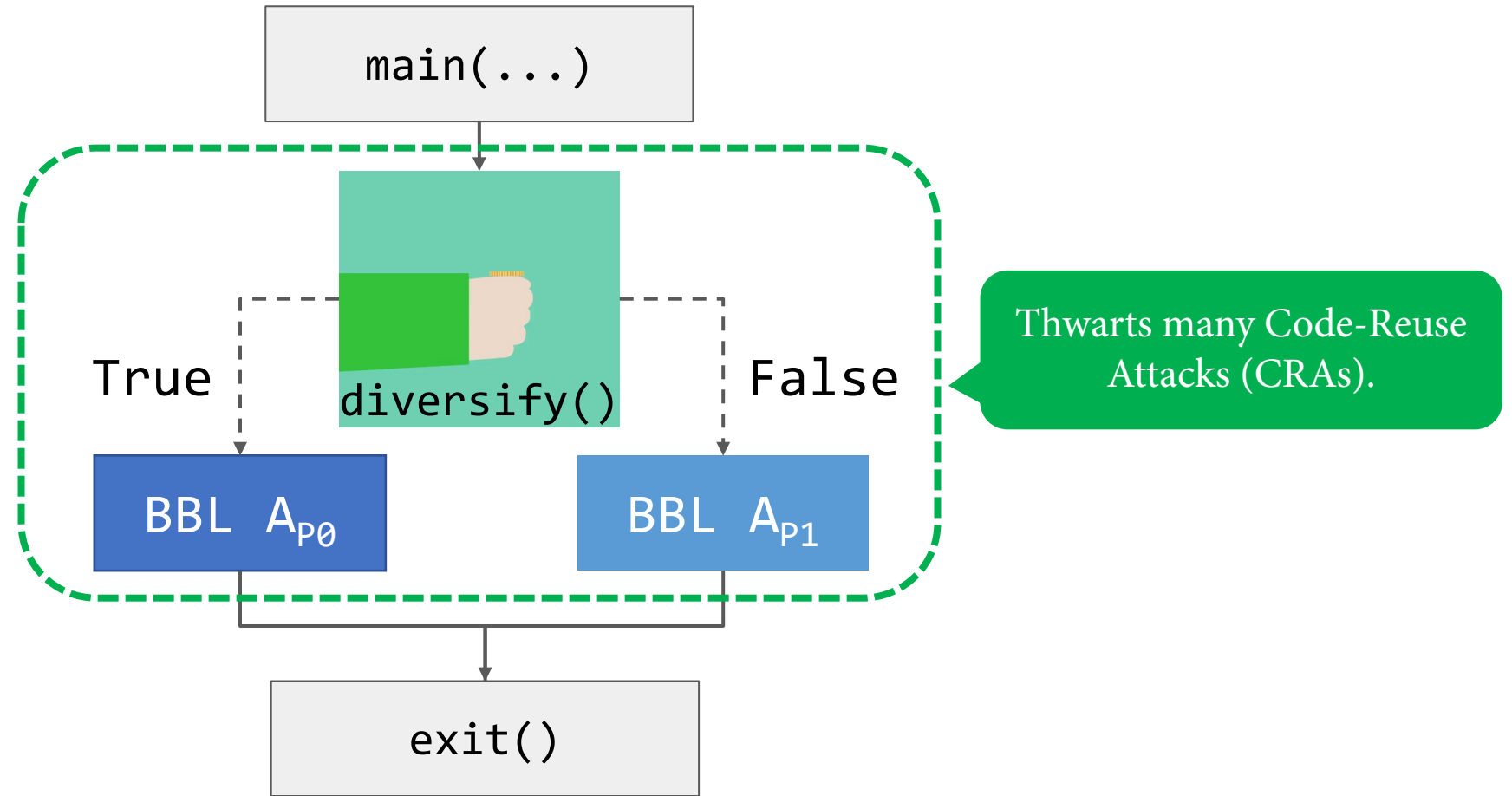
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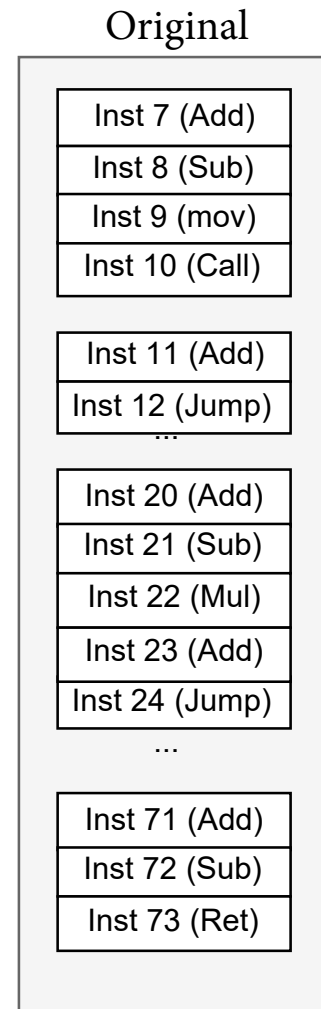
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Program Control Flow Graph

How does PNS protect against CRAs?

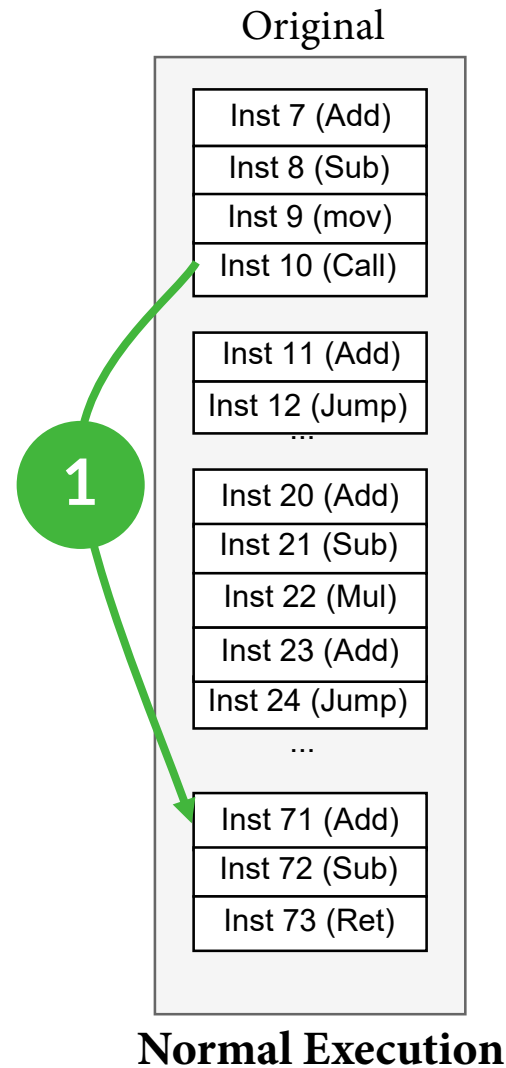
Phantoms force an adversary to guess the execution path.



Normal Execution

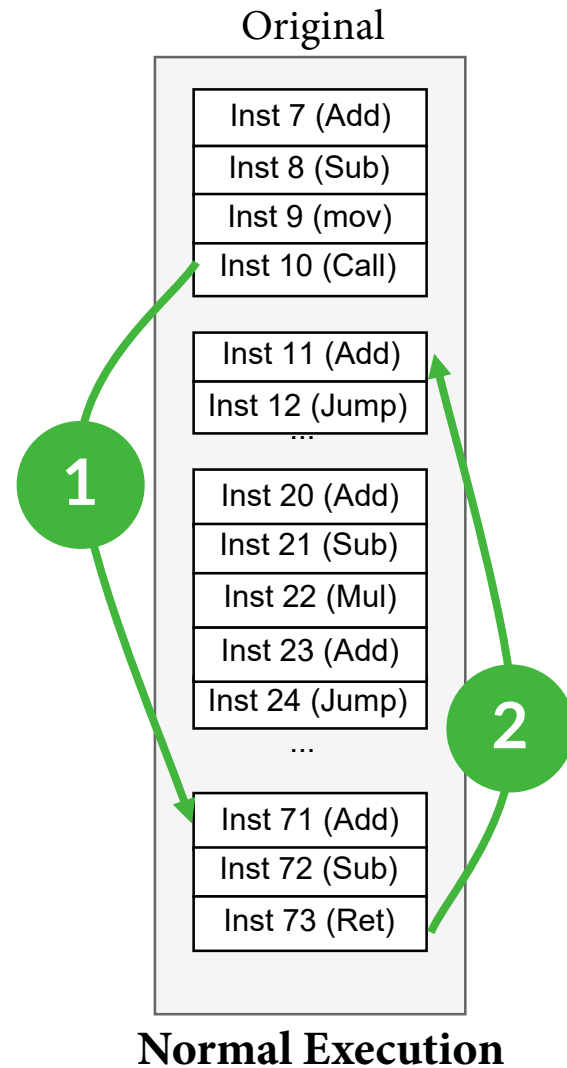
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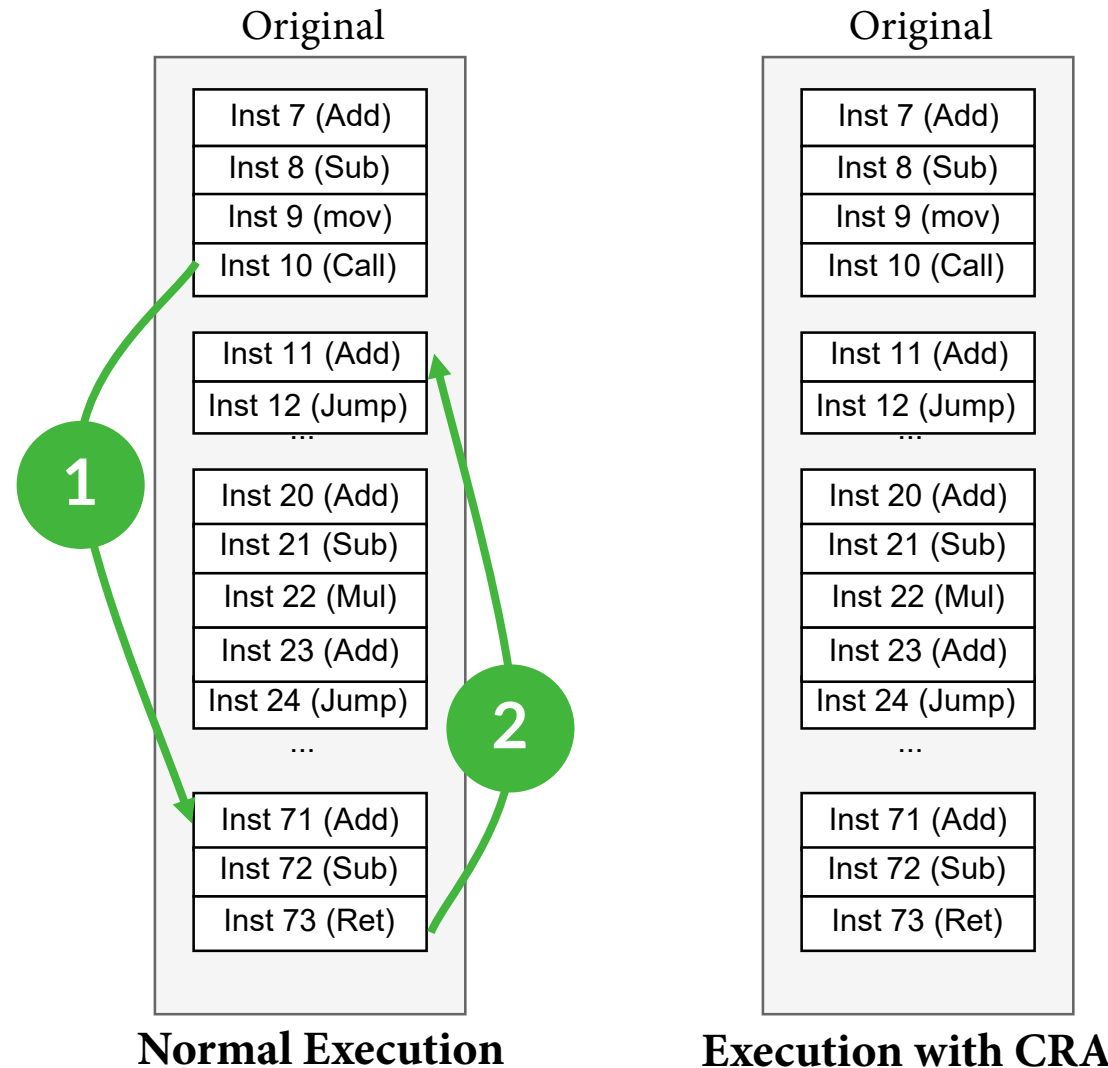
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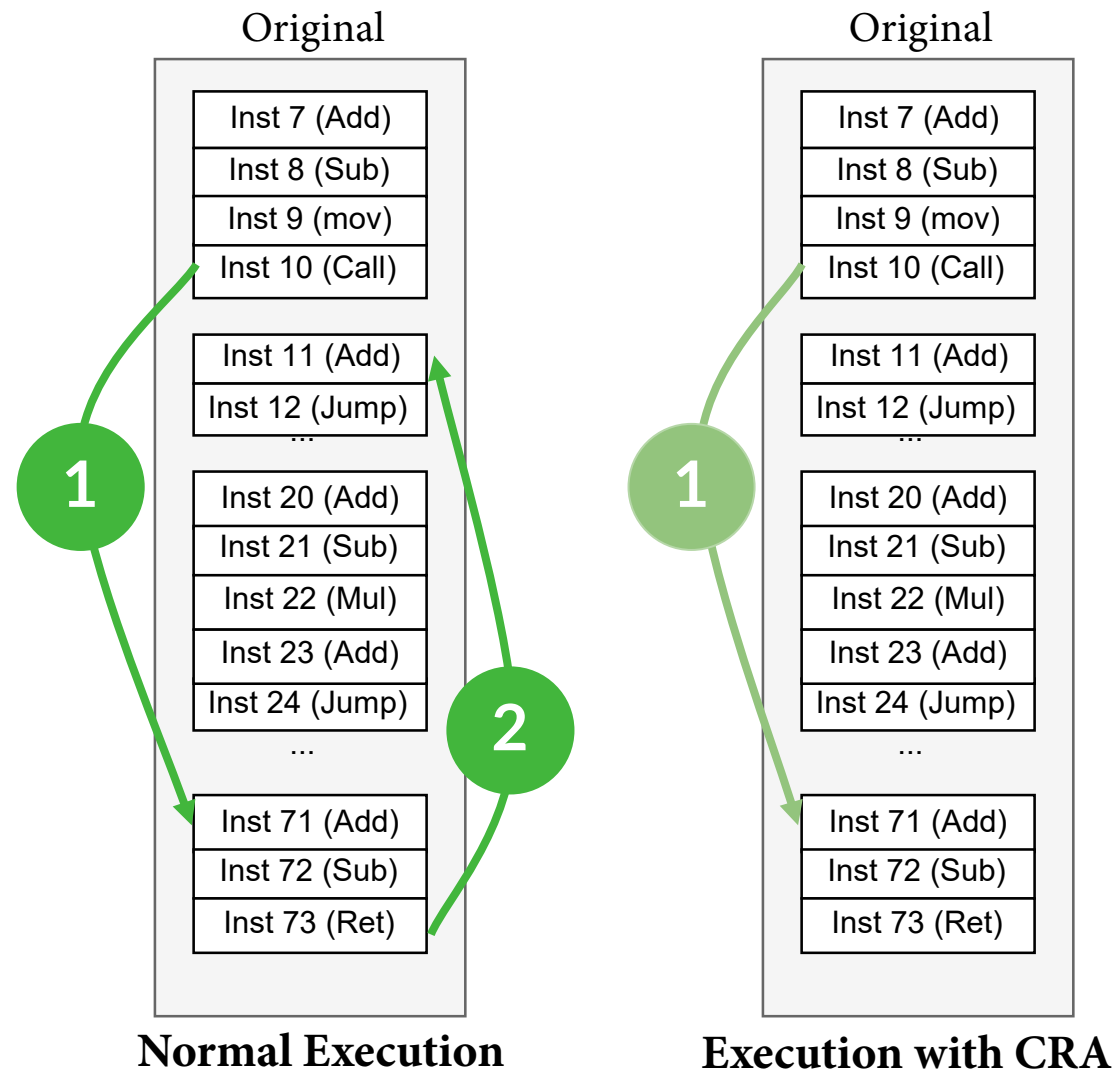
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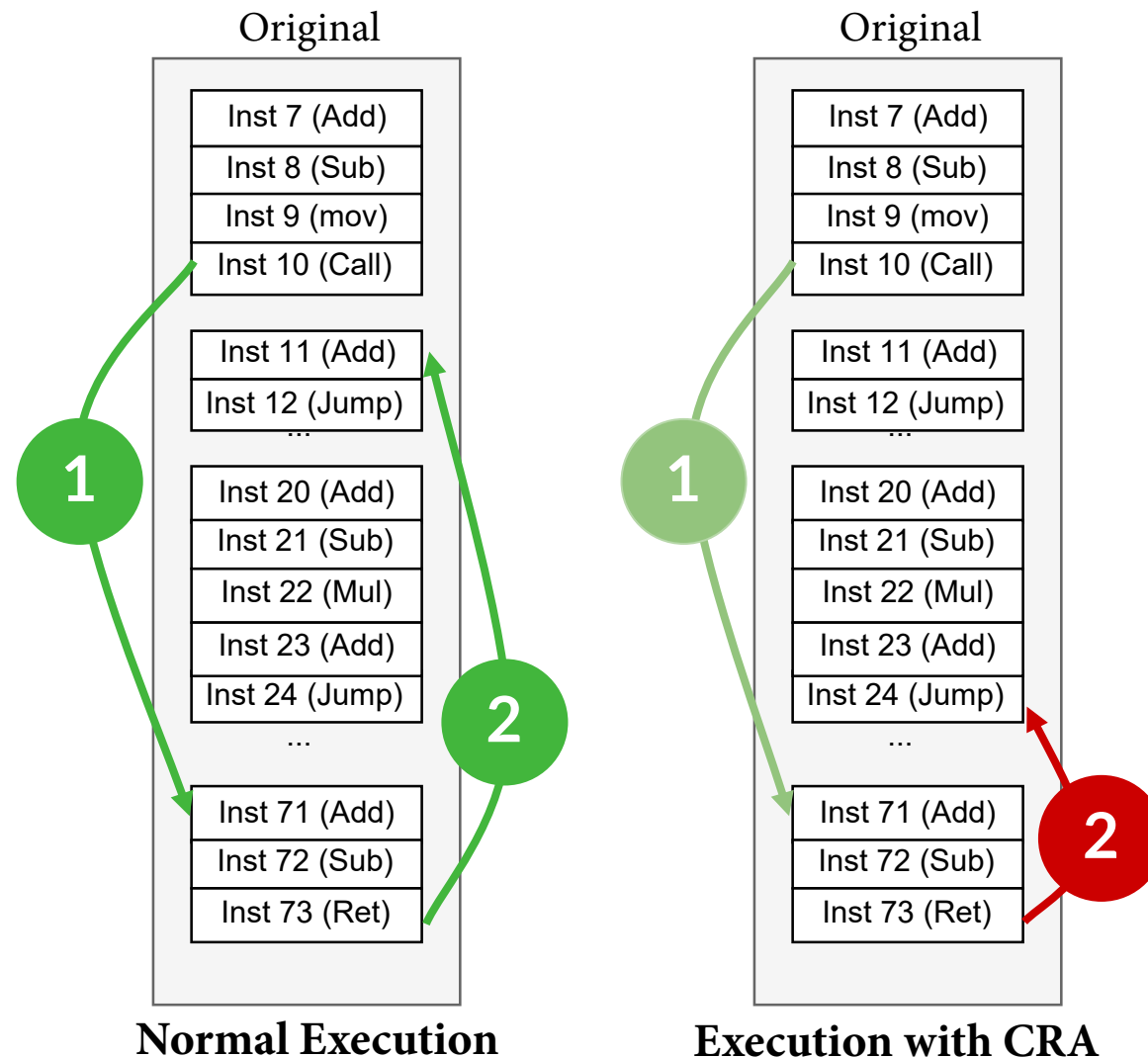
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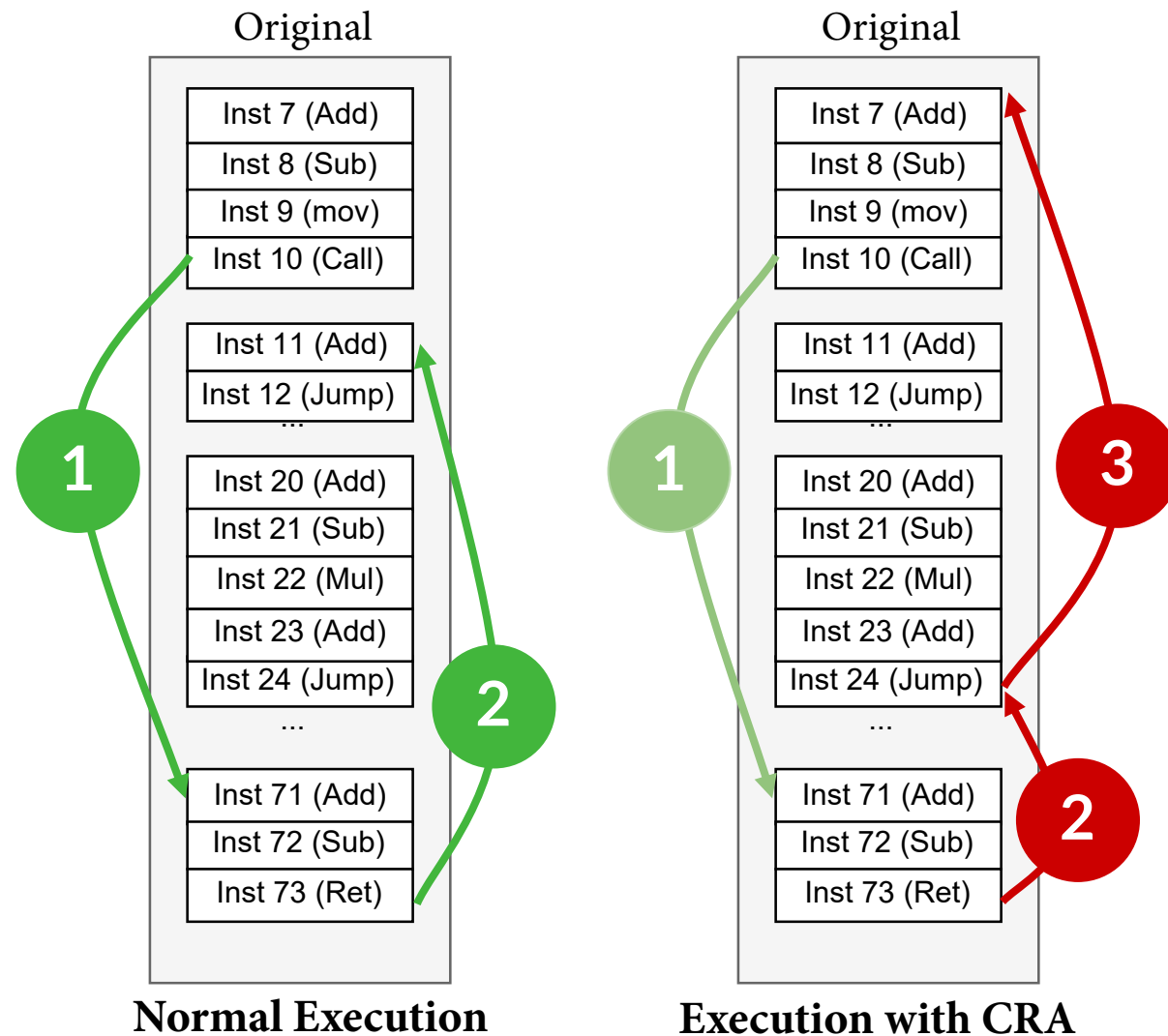
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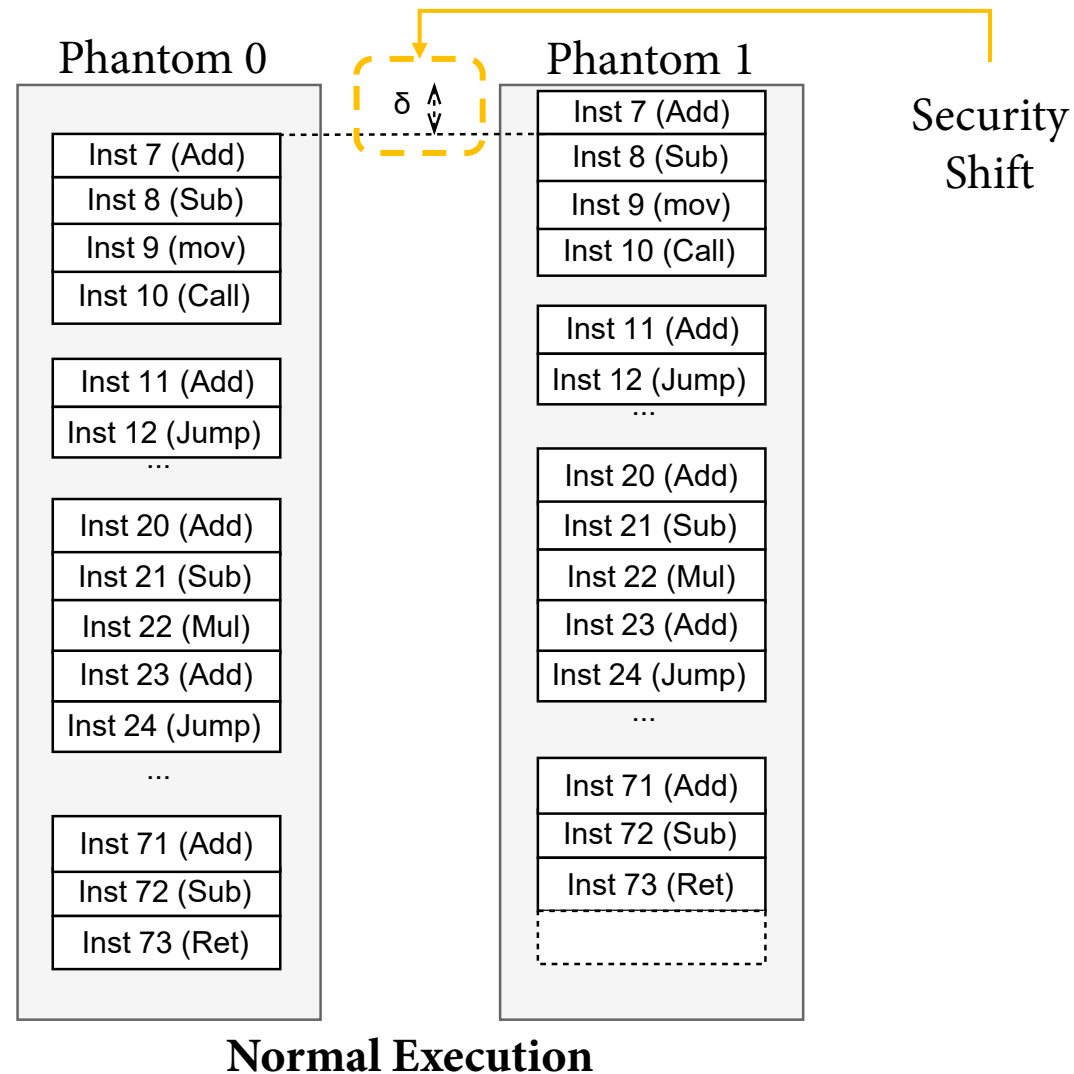
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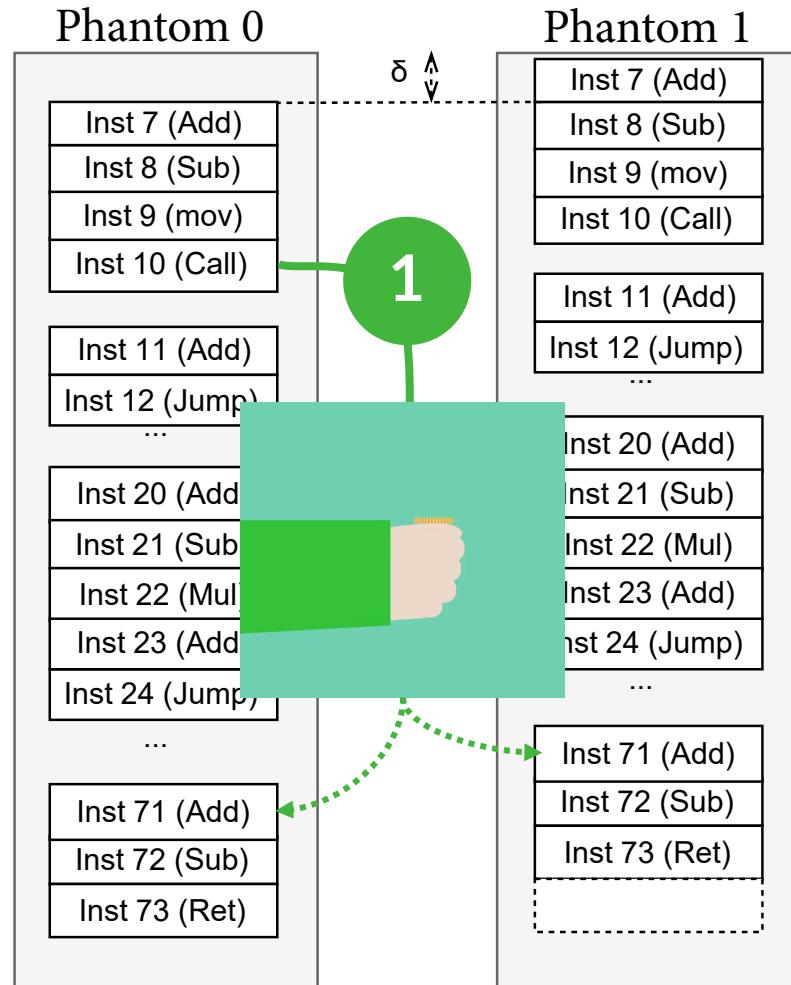
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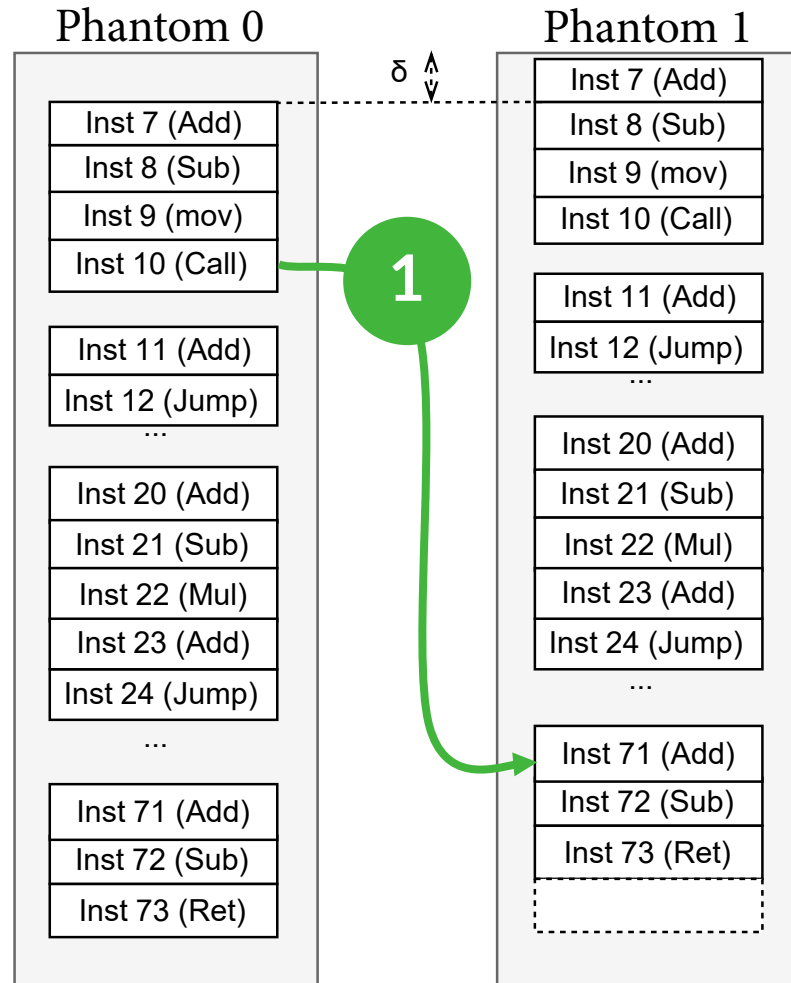
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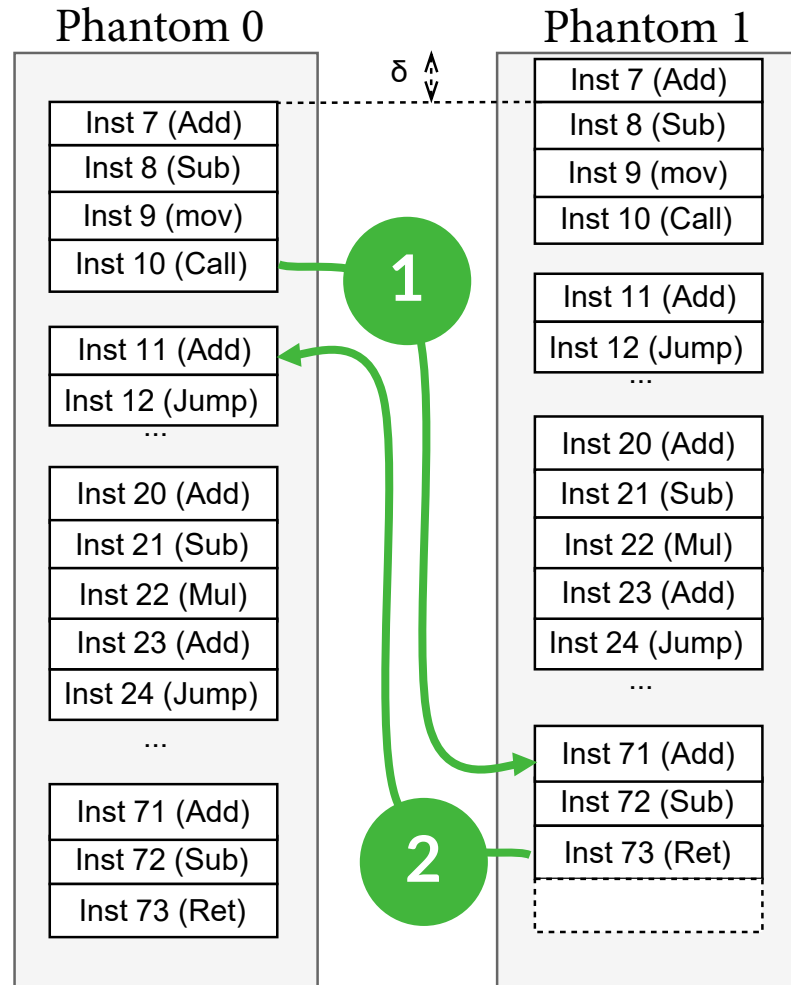
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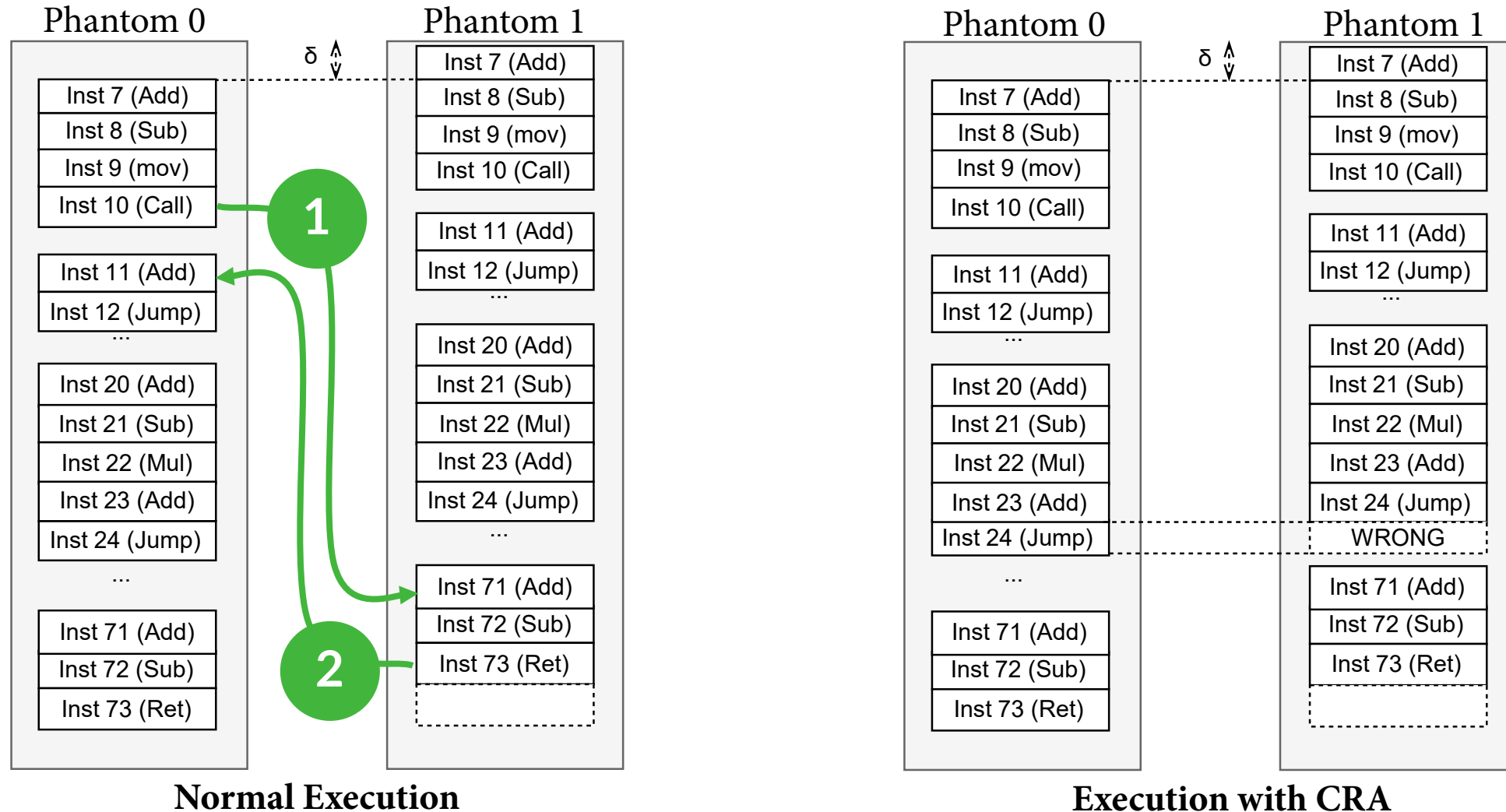
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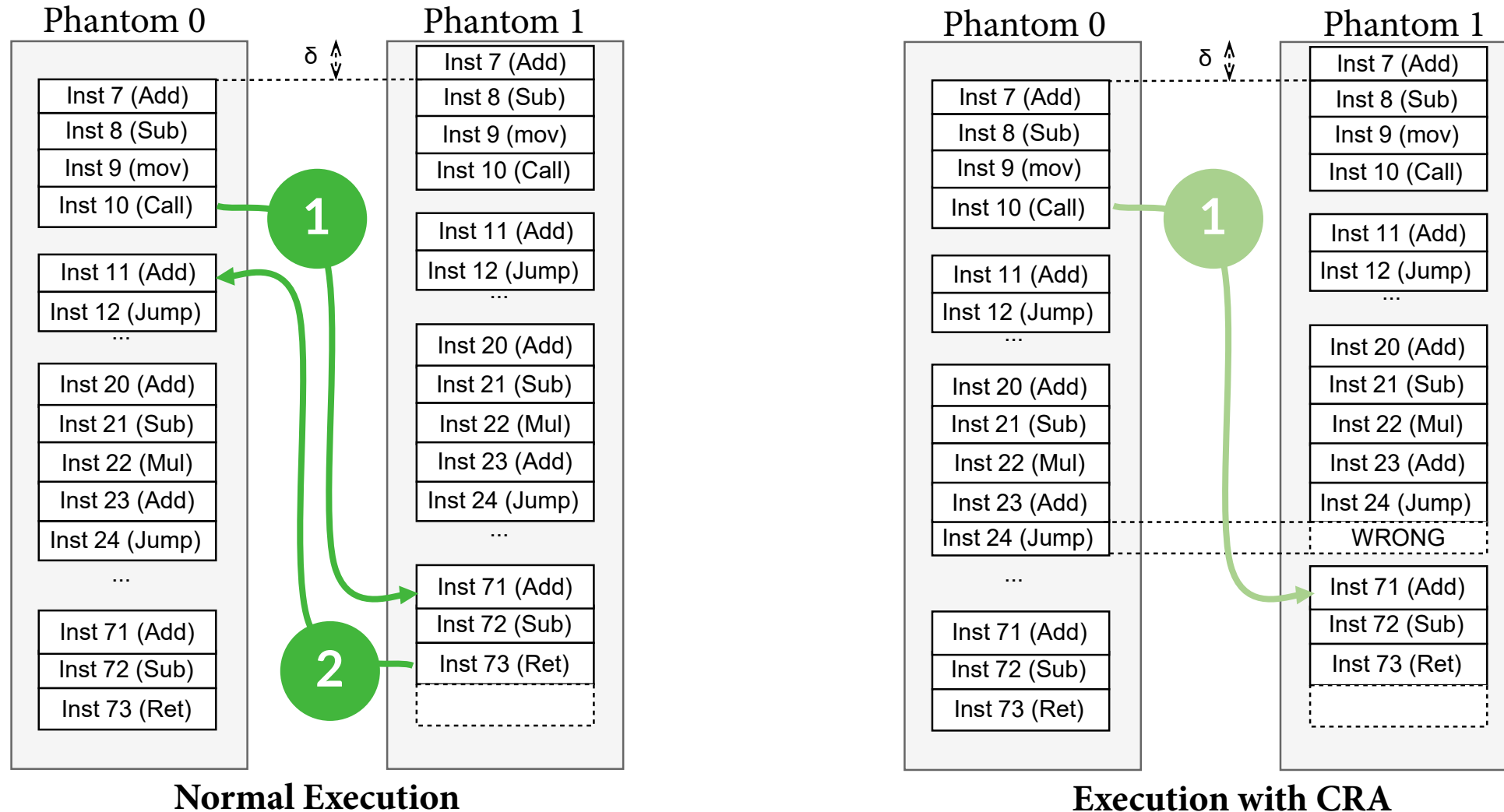
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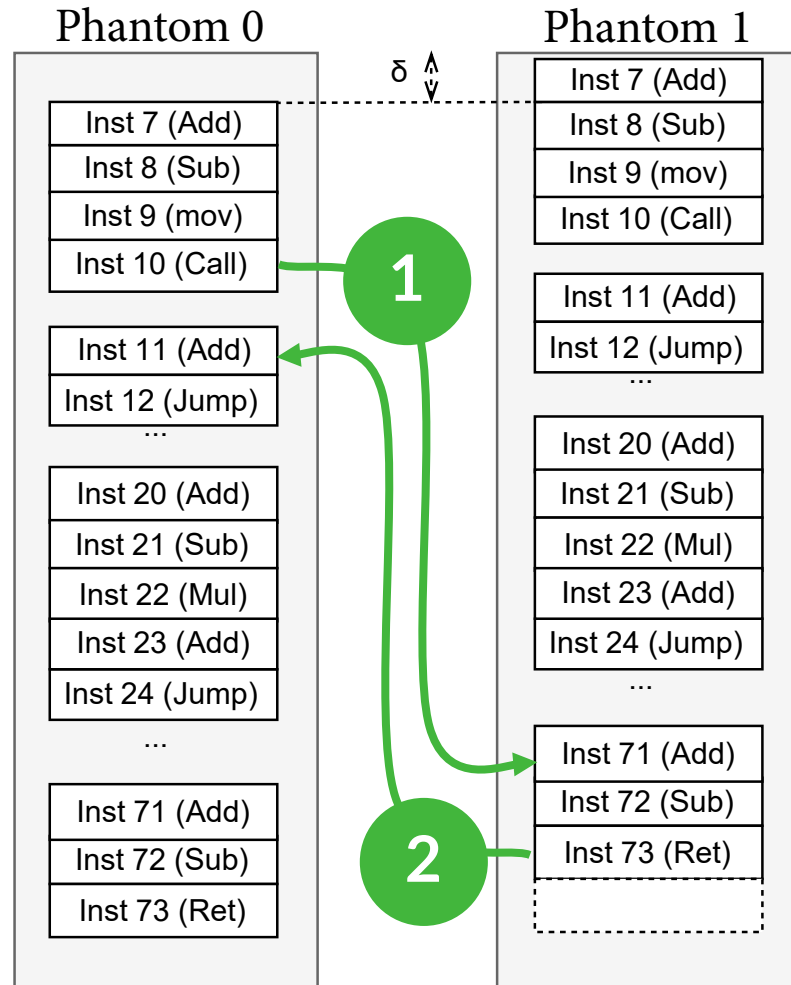
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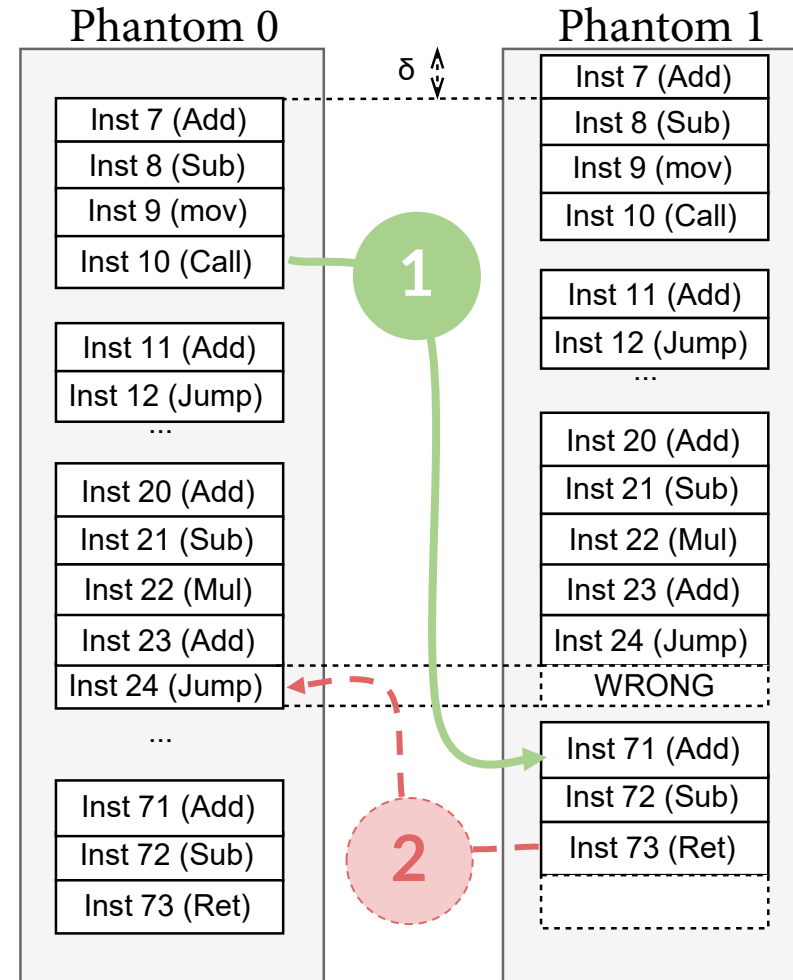


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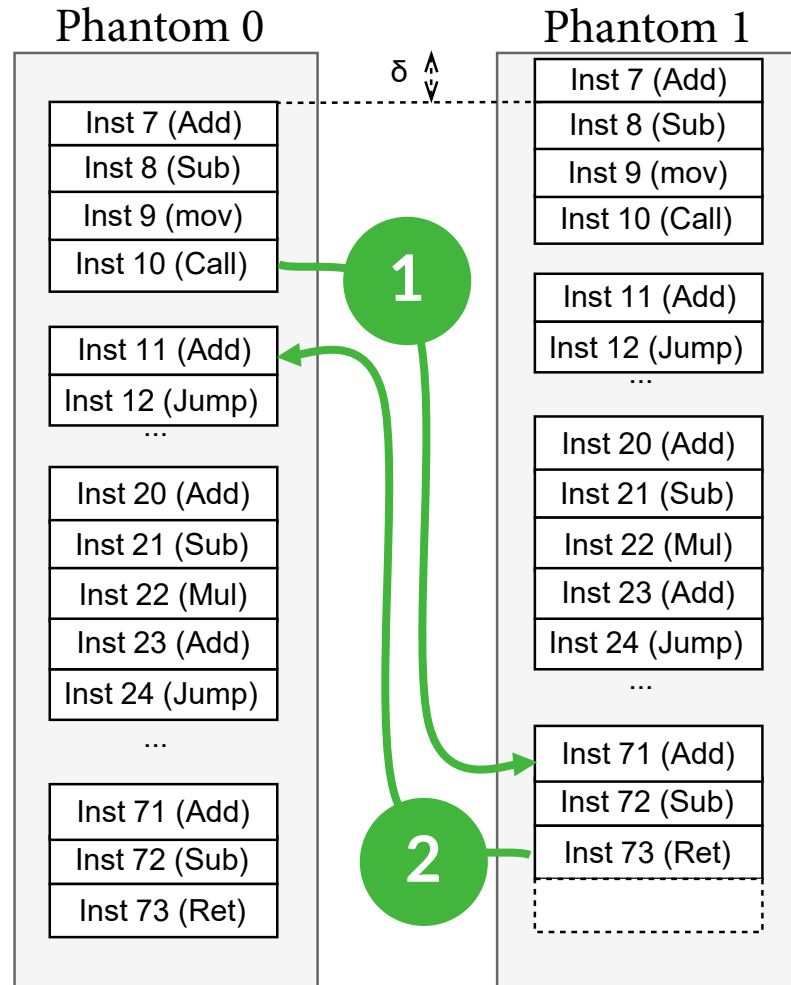
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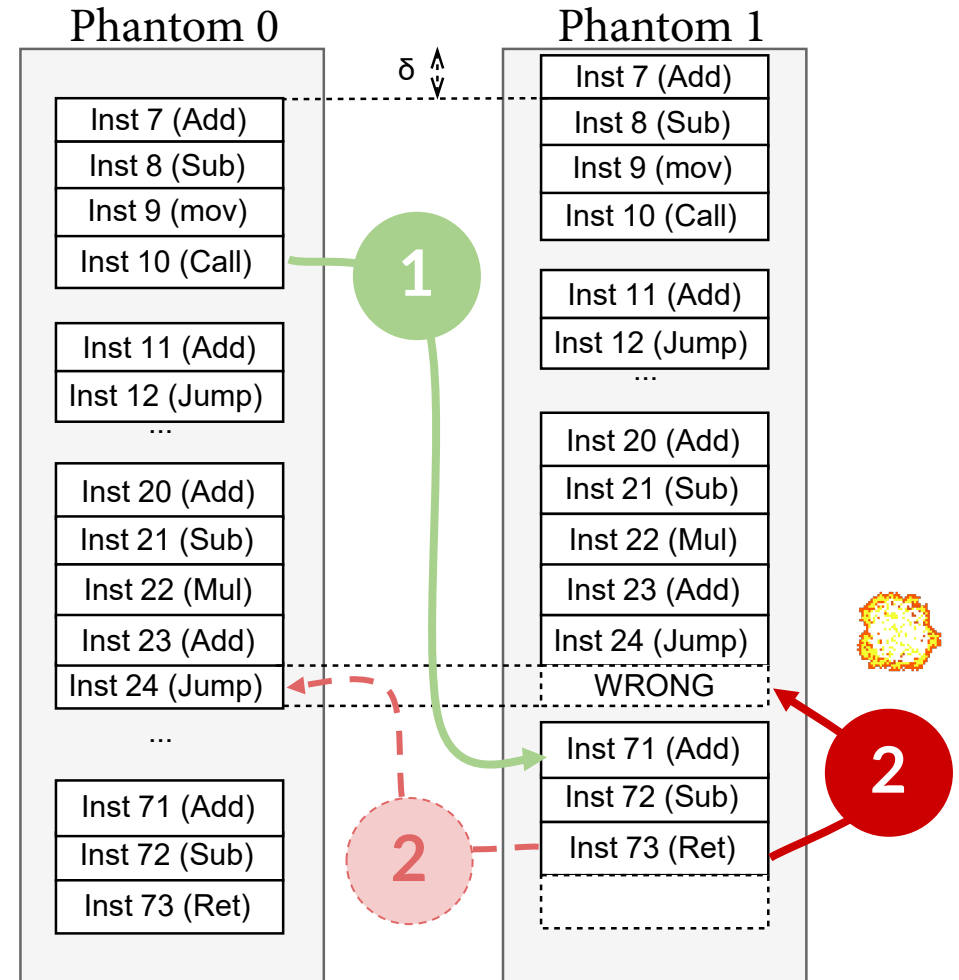
Execution with CRA

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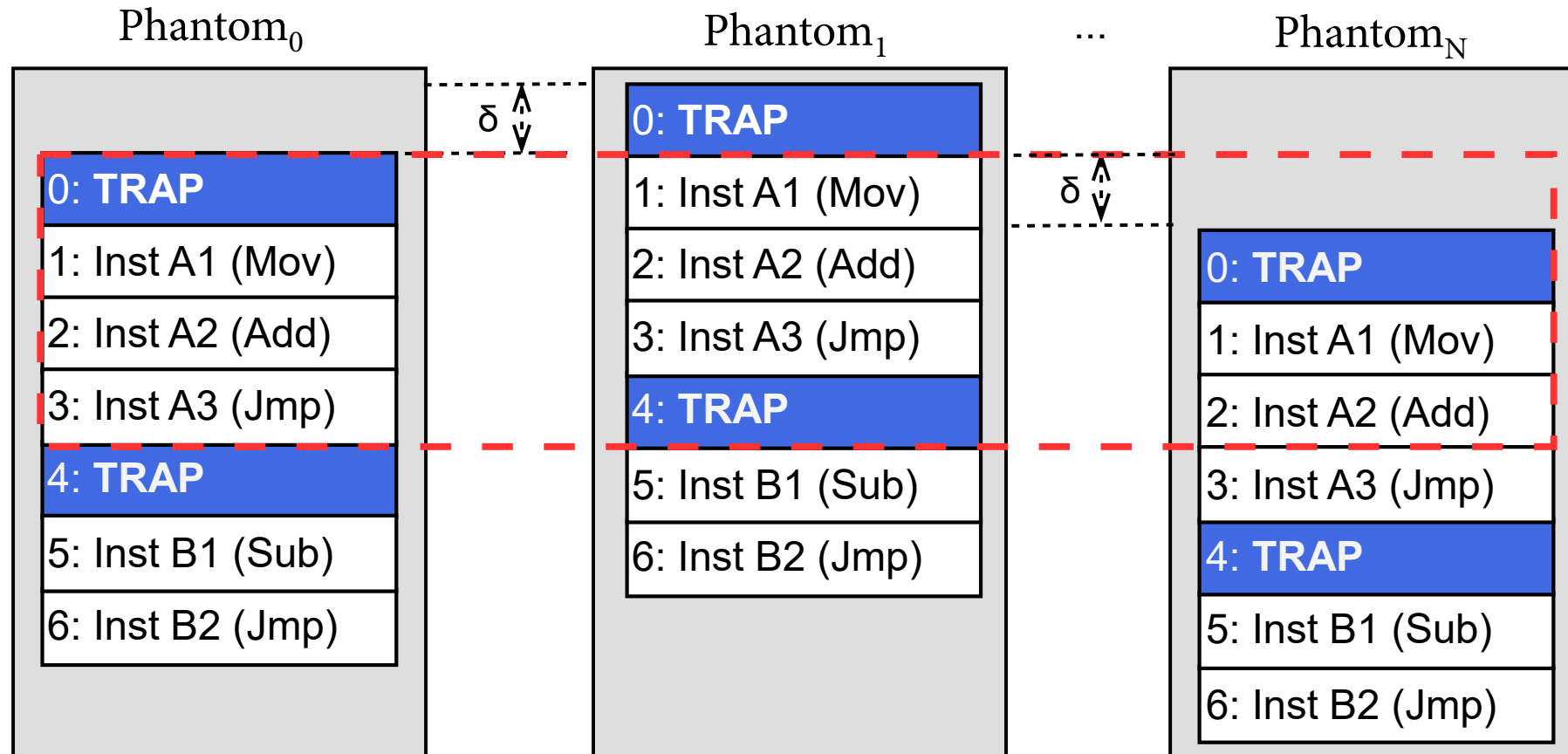
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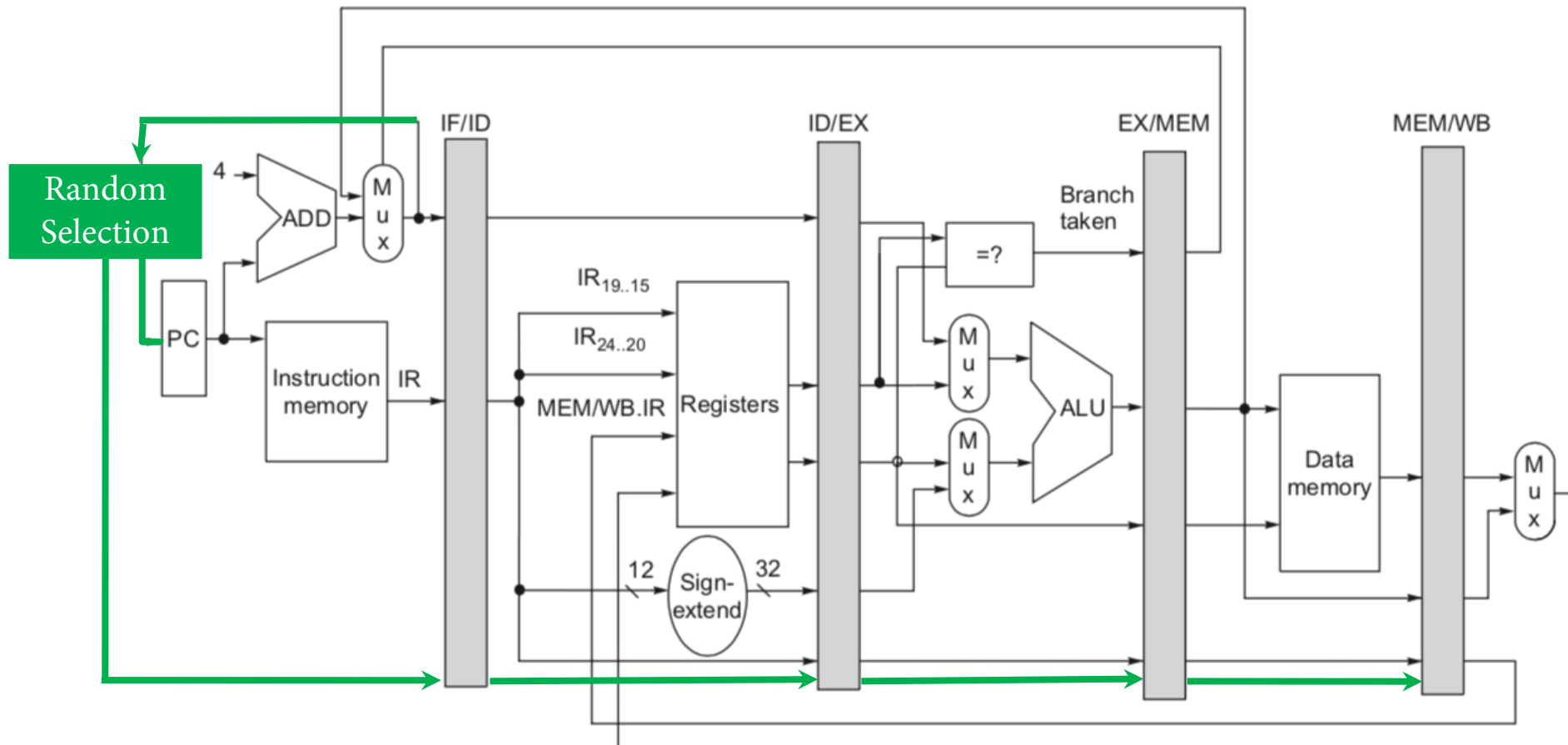
How does PNS precisely trap an attacker?

Code is instrumented with special instructions to throw an exception.



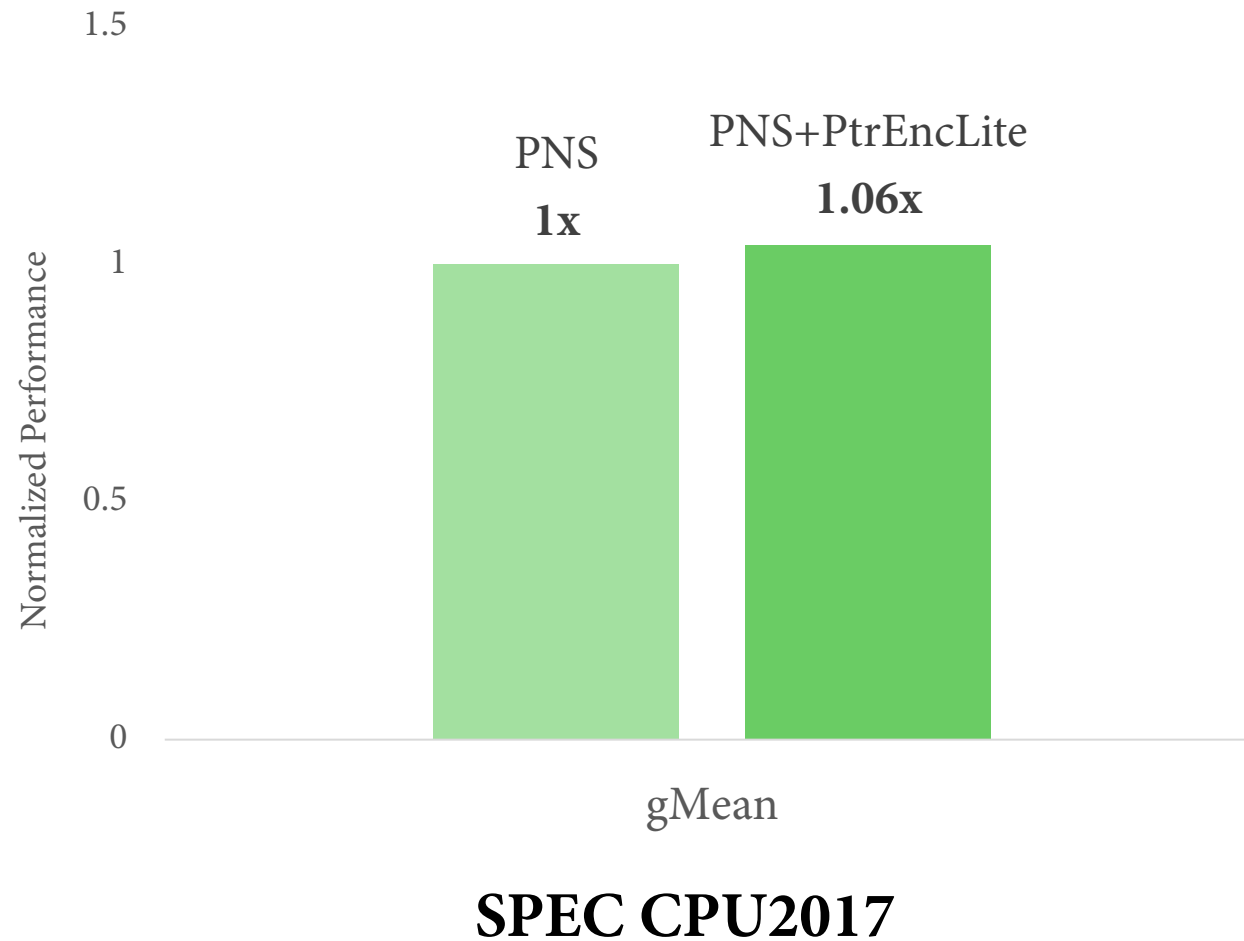
How is PNS implemented?

We do minimal changes to the processor frontend.



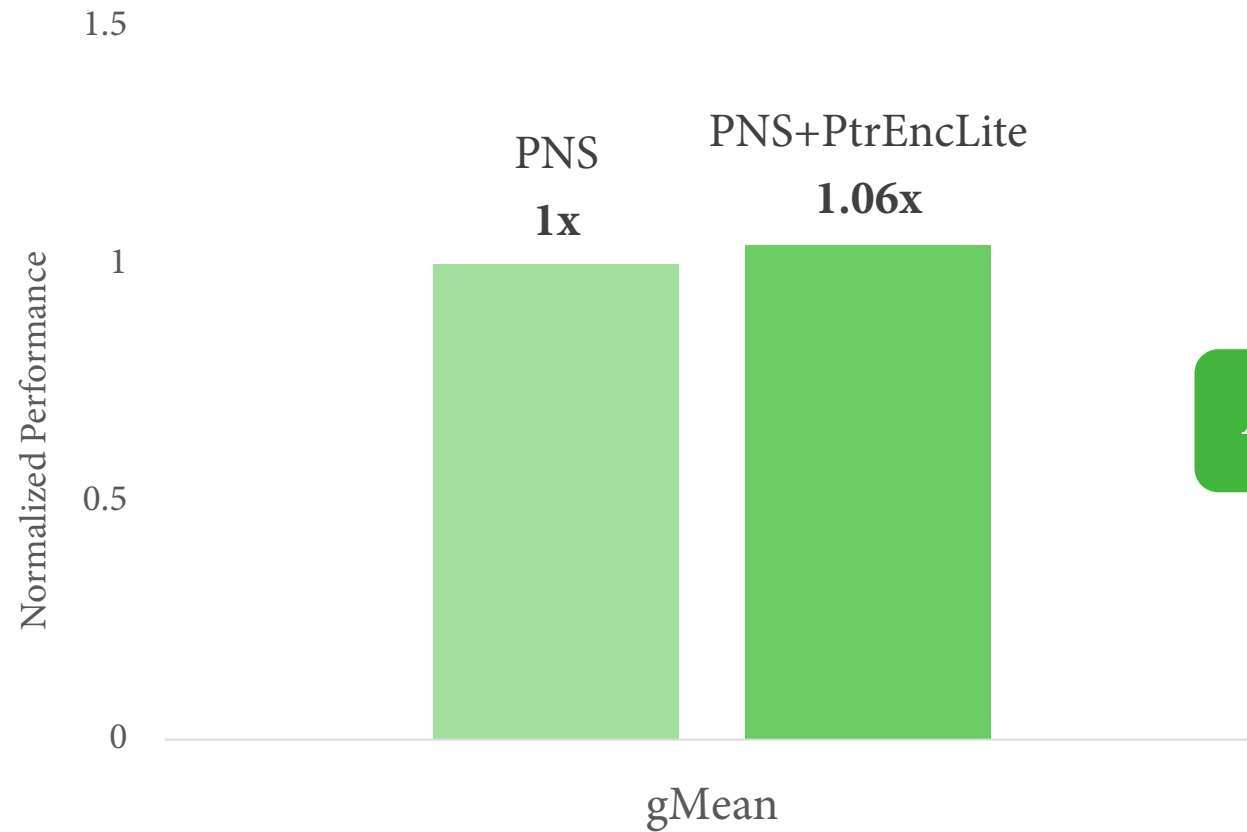
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We used the gem5 architectural simulator to validate correctness & performance.



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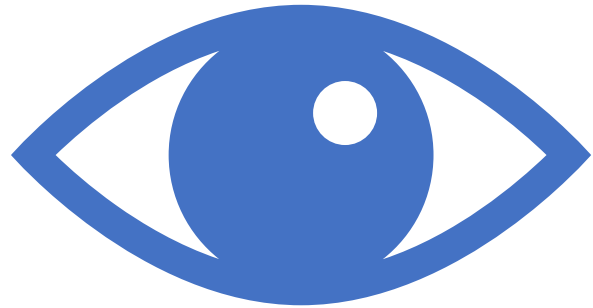


Average slowdown < 6%!

SPEC CPU2017



Limitations



Repeated Observation Attack

Running the same binary on a non-protected system can leak the security shift of a return address.



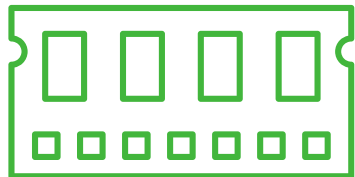
Why is PNS well suited for constrained devices?

It brings efficient protection with minimal cost.



Minimal Performance Impact

PNS has minimal impact on workload execution.

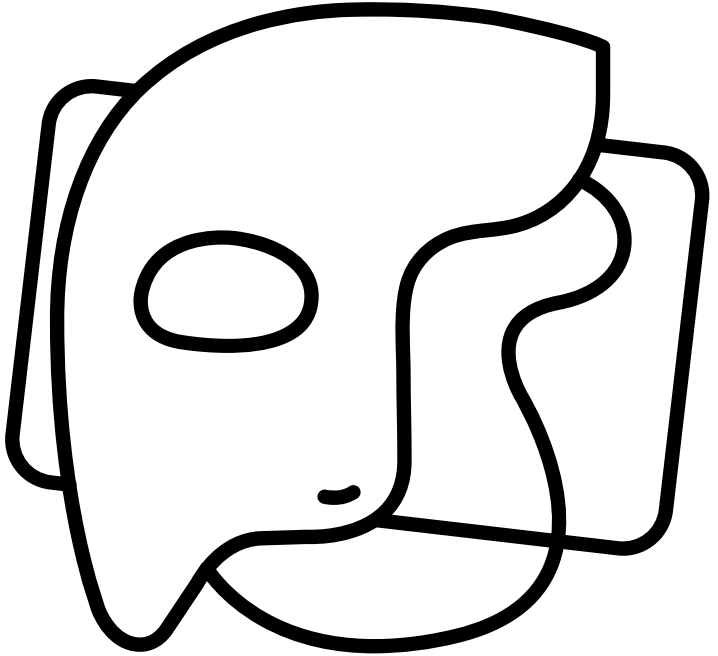


Memory Savings

PNS cuts down on resource duplication associated with aliasing multiple instructions.



Thanks for listening!



Phantom Name System

Find the paper here!



<https://arxiv.org/pdf/1911.02038.pdf>